

# Quantum computation with Turaev-Viro codes

Robert König

joint work with Greg Kuperberg and Ben Reichardt

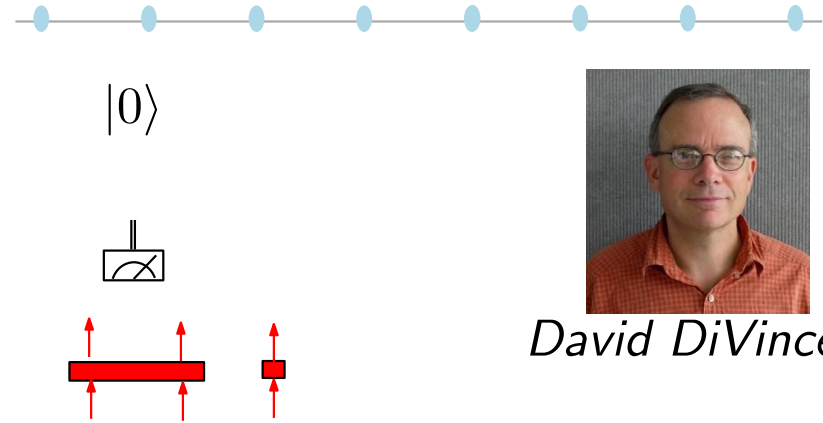
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  - mapping class group representations
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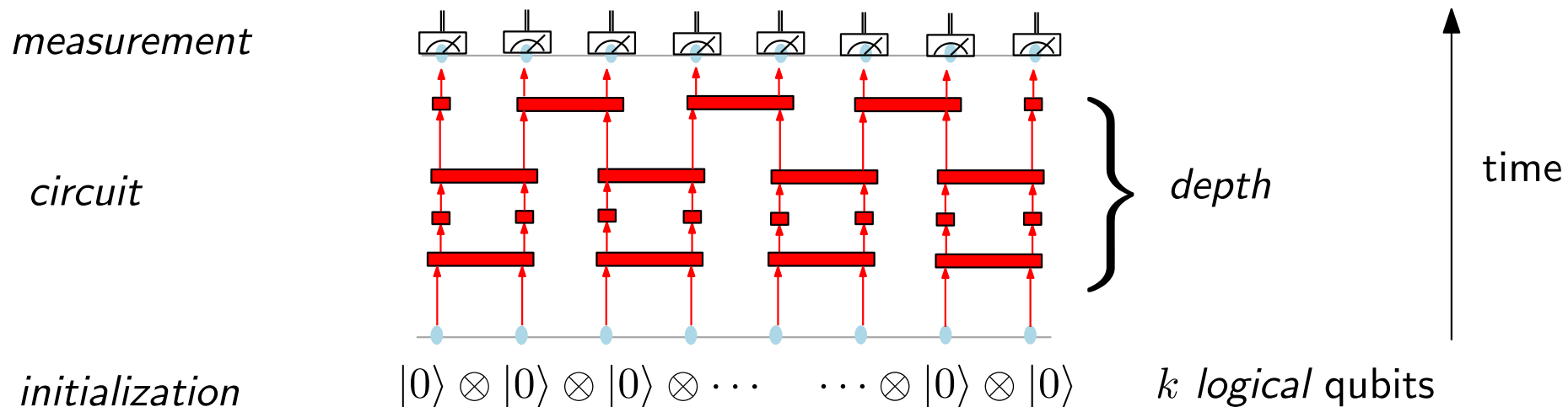
# Quantum fault-tolerance: the DiVincenzo criteria

## DiVincenzo criteria for fault-tolerant quantum computation

1. scalable physical system with well-characterized qubits
2. ability to initialize fiducial state
3. decoherence times  $\gg$  gate operation time
4. qubit-specific measurement capability
5. universal set of quantum gates



David DiVincenzo



# Quantum noise on n qubits

**Quantum noise** on n qubits is represented by a completely positive trace-preserving map (CPTPM)

$$\mathcal{N} : \mathcal{B}((\mathbb{C}^2)^{\otimes n}) \rightarrow \mathcal{B}((\mathbb{C}^2)^{\otimes n})$$

**Operational problem:** can we recover information subjected to such noise?

Using the Kraus decomposition  $\mathcal{N}(\rho) = \sum_{E \in \mathcal{E}} E \rho E^\dagger$

it can be shown that it suffices to protect against a certain set of errors  $\mathcal{E}$

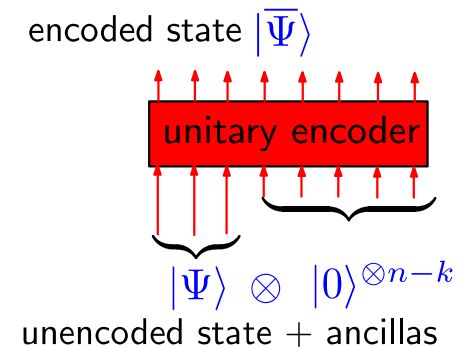
where an error is a linear map  $E : (\mathbb{C}^2)^{\otimes n} \rightarrow (\mathbb{C}^2)^{\otimes n}$

**Mathematical problem:** Is there a recovery CPTPM  $\mathcal{R} : \mathcal{B}((\mathbb{C}^2)^{\otimes n}) \rightarrow \mathcal{B}((\mathbb{C}^2)^{\otimes n})$

such that for "suitable"  $\rho$   $\mathcal{R}(E \rho E^\dagger) \propto \rho$  for all  $E \in \mathcal{E}$

**Procedure:** (isometrically) embed/ "encode"

$$\begin{aligned} (\mathbb{C}^2)^{\otimes k} &\rightarrow \mathcal{L} \subset (\mathbb{C}^2)^{\otimes n} \\ \Psi &\mapsto \bar{\Psi} \end{aligned}$$



**QEC condition:**  
[Knill, Laflamme]

$$\mathcal{L} \text{ protects against errors } \mathcal{E} \iff \langle \bar{\Psi} | E^\dagger F | \bar{\varphi} \rangle = c(E, F) \langle \bar{\Psi} | \bar{\varphi} \rangle$$

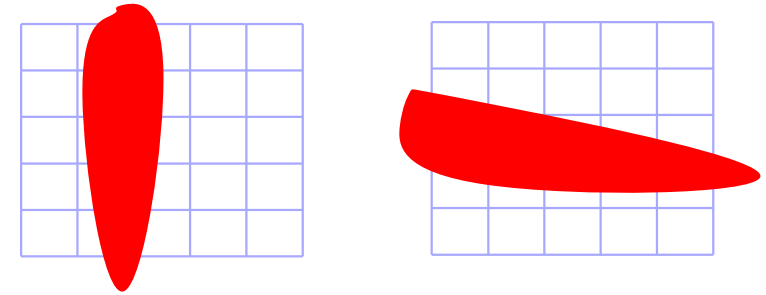
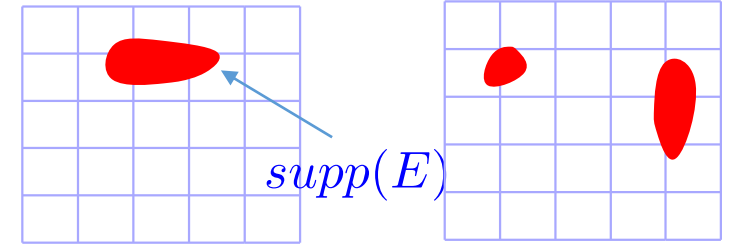
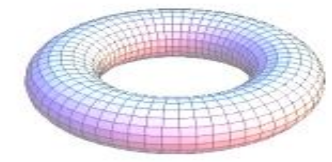
for all  $E, F \in \mathcal{E}, \bar{\Psi}, \bar{\varphi} \in \mathcal{L}$

# “Topological” error-correcting codes

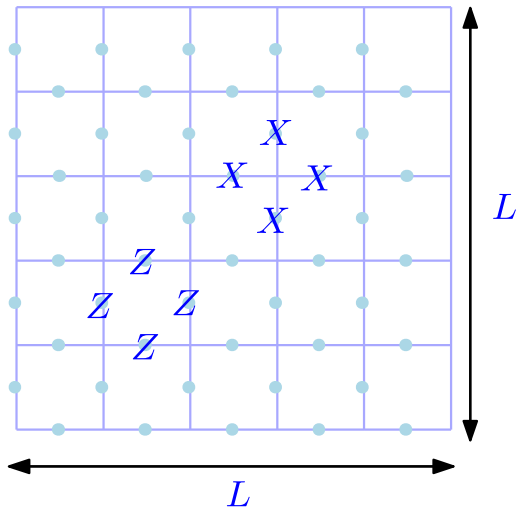
**Def:** A “topological” code:

protects against all local errors, e.g.,  
and more generally errors with  
“topologically trivial” support

does not protect against errors  
with topologically non-trivial  
support, e.g.,



## Example: Kitaev’s toric code



$n = 2L^2$  qubits on the edges of a edges of a  $L \times L$  periodic lattices

$$\mathcal{L} = \{ \Psi \in (\mathbb{C}^2)^{\otimes n} \mid A_v \Psi = B_p \Psi = \Psi \quad \text{for all } v, p \}$$

$$A_v = X^{\otimes 4} \text{ for each vertex } v$$

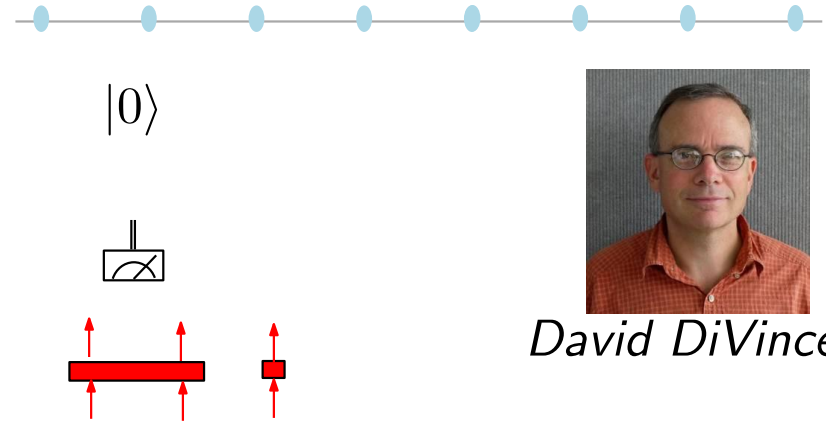
$$B_p = Z^{\otimes 4} \text{ for each plaquette } p$$

$$k = \log_2 \dim \mathcal{L} = 2 \text{ encoded qubits}$$

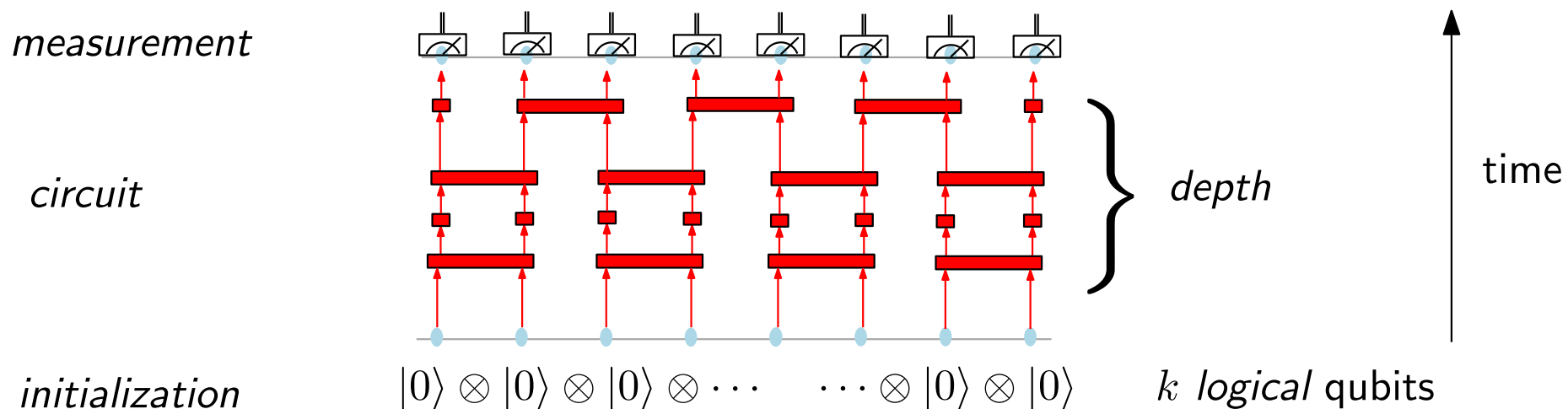
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- ✓ 3. decoherence times  $\gg$  gate operation time
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- ➔ 5. universal set of quantum gates



David DiVincenzo



The code space of Kitaev's toric code

# Logical operators in Kitaev's toric code

The operators  $\bar{X}_1, \bar{Z}_1, \bar{X}_2, \bar{Z}_2$

- preserve the code space  $\mathcal{L}$ , i.e., are *logical*
- satisfy Pauli commutation relations

$\Rightarrow$  They define a factorization of the code space  $\mathcal{L} \cong \mathbb{C}^2 \otimes \mathbb{C}^2$  such that

$$\begin{aligned} \bar{X}_1 &\cong X \otimes I \\ \bar{Z}_1 &\cong Z \otimes I \\ \bar{X}_2 &\cong X \otimes I \\ \bar{Z}_2 &\cong I \otimes X \end{aligned}$$

$$\bar{X}_1 = \begin{array}{|c|c|c|c|c|} \hline & & X & & \\ \hline & & X & & \\ \hline & & X & & \\ \hline & & X & & \\ \hline & & X & & \\ \hline & & X & & \\ \hline & & X & & \\ \hline \end{array}$$

$$\bar{Z}_2 = \begin{array}{|c|c|c|c|c|} \hline & & Z & & \\ \hline & & Z & & \\ \hline & & Z & & \\ \hline & & Z & & \\ \hline & & Z & & \\ \hline & & Z & & \\ \hline \end{array}$$

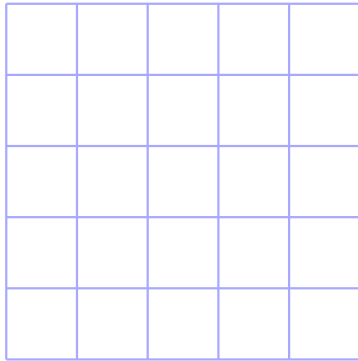
$$\bar{X}_2 = \begin{array}{|c|c|c|c|c|c|} \hline & & & & & \\ \hline & & & & & \\ \hline & & & & & \\ \hline X & X & X & X & X & X \\ \hline & & & & & \\ \hline & & & & & \\ \hline \end{array}$$

$$\bar{Z}_1 = \begin{array}{|c|c|c|c|c|c|} \hline & & & & & \\ \hline & & & & & \\ \hline & & & & & \\ \hline Z & Z & Z & Z & Z & \\ \hline & & & & & \\ \hline & & & & & \\ \hline \end{array}$$

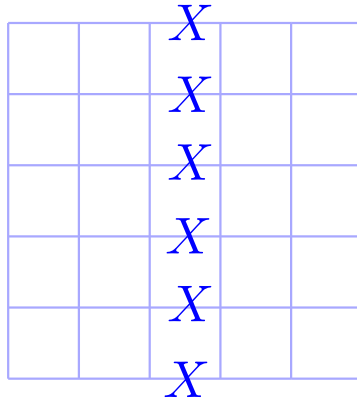


# Logical operators in Kitaev's toric code: commuting subalgebras

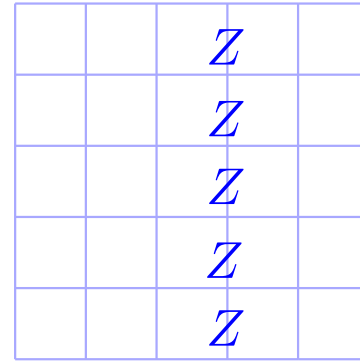
$F_{(0,0)}(C_1)$



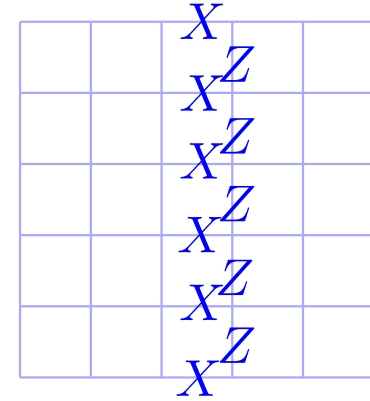
$F_{(1,0)}(C_1)$



$F_{(0,1)}(C_1)$

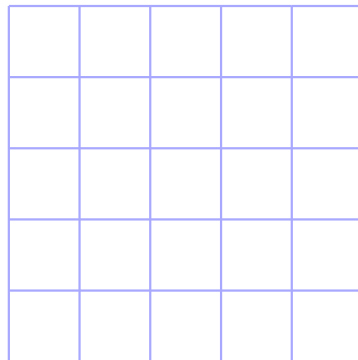


$F_{(1,1)}(C_1)$

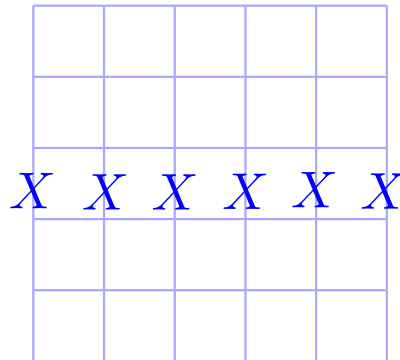


these 4  
commute:

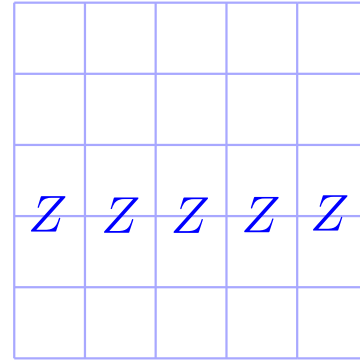
$F_{(0,0)}(C_2)$



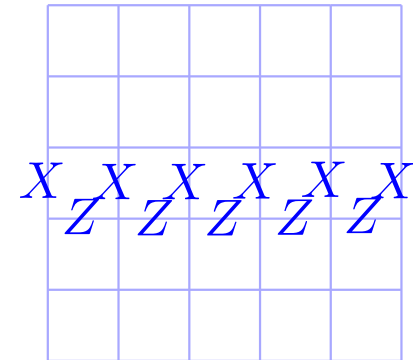
$F_{(1,0)}(C_2)$



$F_{(0,1)}(C_2)$

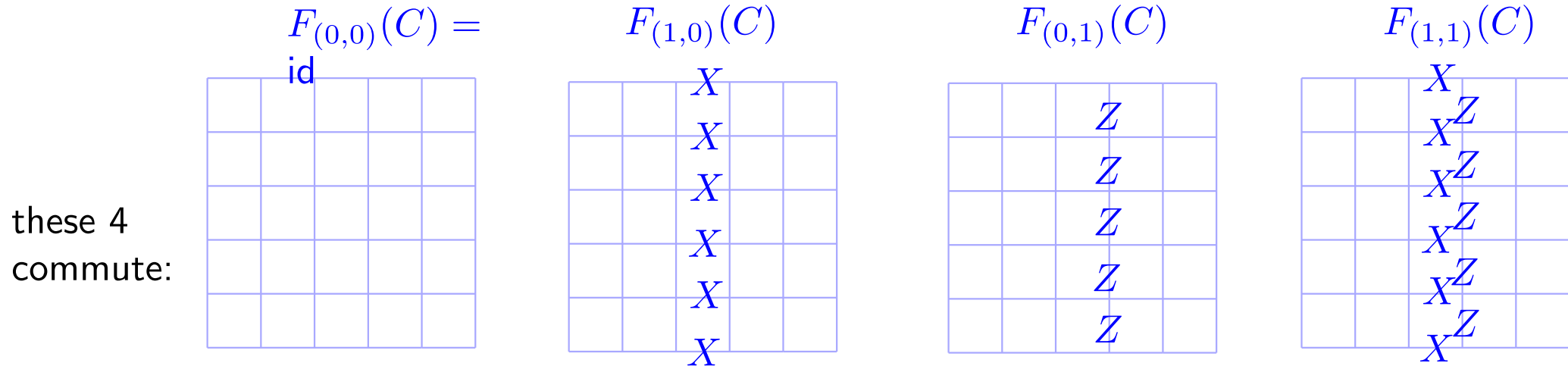


$F_{(1,1)}(C_2)$



these 4  
commute:

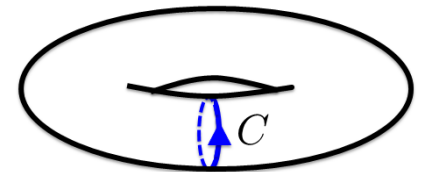
# ``Flux''-basis states associated with loops on a torus



we can use the following 4 orthogonal projections to label basis states of the code space:

(These are the idempotents of the **Verlinde algebra**.)

$$\begin{aligned}
 P_{(0,0)}(C) &= \frac{1}{2}(\text{id} + X^{\otimes L}) \cdot \frac{1}{2}(\text{id} + Z^{\otimes L}) & |1\rangle_C \\
 P_{(1,0)}(C) &= \frac{1}{2}(\text{id} - X^{\otimes L}) \cdot \frac{1}{2}(\text{id} + Z^{\otimes L}) & |e\rangle_C \\
 P_{(0,1)}(C) &= \frac{1}{2}(\text{id} + X^{\otimes L}) \cdot \frac{1}{2}(\text{id} - Z^{\otimes L}) & |m\rangle_C \\
 P_{(1,1)}(C) &= \frac{1}{2}(\text{id} - X^{\otimes L}) \cdot \frac{1}{2}(\text{id} - Z^{\otimes L}) & |\epsilon\rangle_C
 \end{aligned}$$



Every non-contractible closed loop  $C$  gives rise to a basis  $\mathcal{B}_C$  of the code space

Fault-tolerant gates (on Kitaev's toric code)

# Fault-tolerant execution logical gates: three ways

## 1) Apply a string-operator

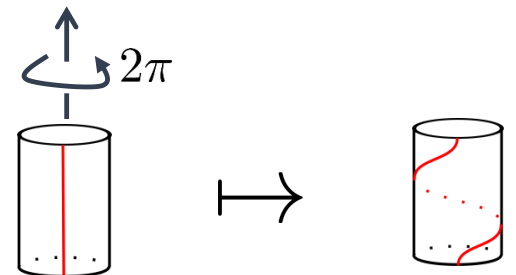
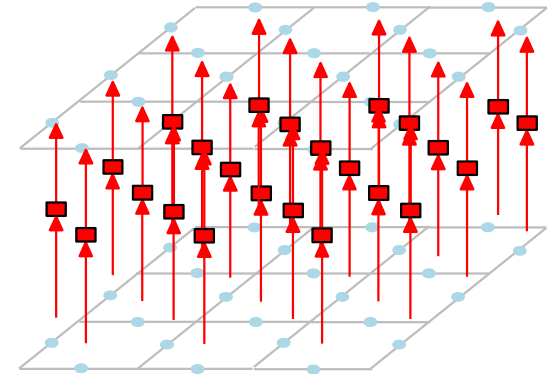
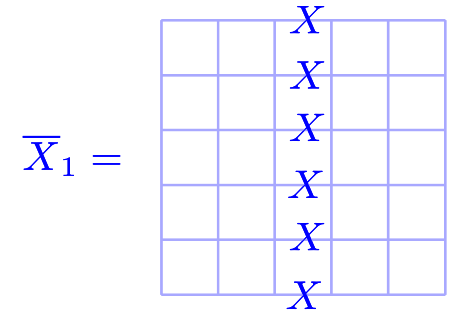
- only gives logical Pauli operators
- does not generalize

## 2) Apply a short (transversal) quantum circuit

- gives certain Clifford operations
- generalization?

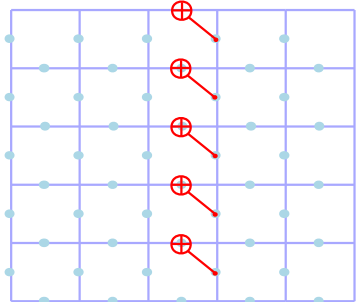
## 3) Apply code deformation (sequence of codes)

- generalizes to other models: mapping class group representation
- gives universal gate sets (in certain models)!



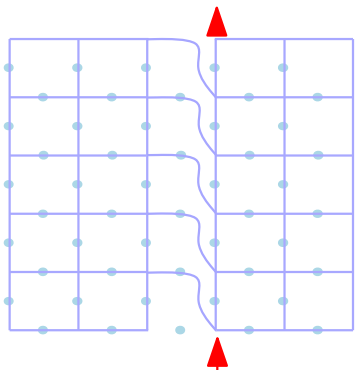
# Mapping class group representation and toric code

apply  $L$  CNOTs  
in parallel

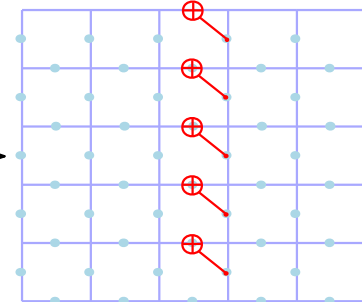


 CNOT gate

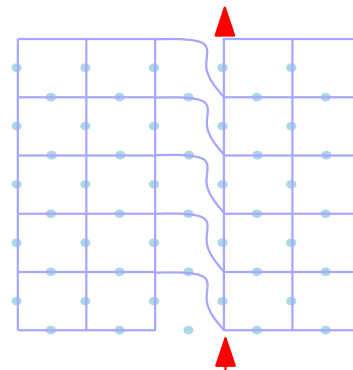
relocate qubits



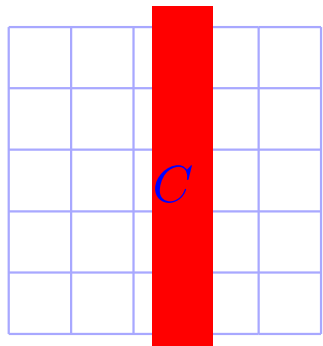
apply  $L$  CNOTs  
in parallel



relocate qubits



repeat  $L$  times

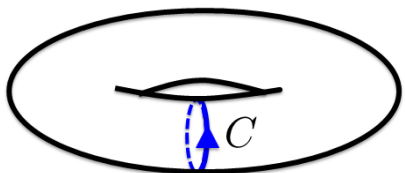
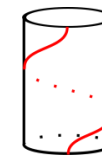
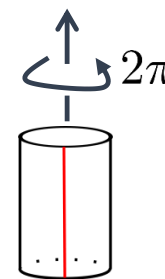


basis states of the code space:

$|1\rangle_C$   
 $|e\rangle_C$   
 $|m\rangle_C$   
 $|\epsilon\rangle_C$

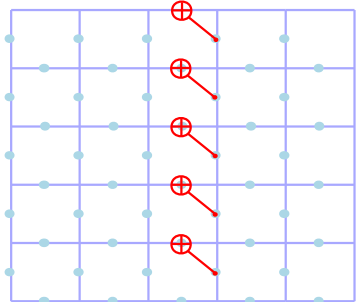
are **eigenvectors**  
**of this operation**  
with eigenvalues

1  
-1  
-1  
1



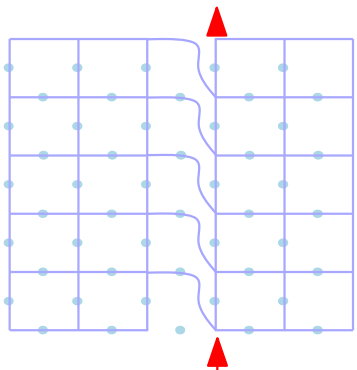
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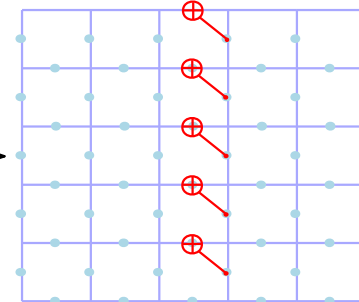


 CNOT gate

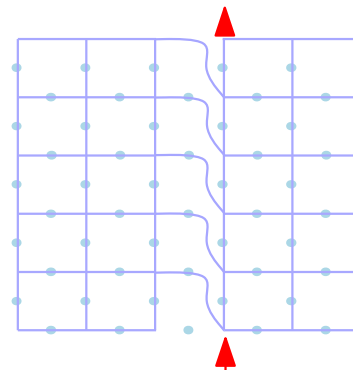
relocate qubits



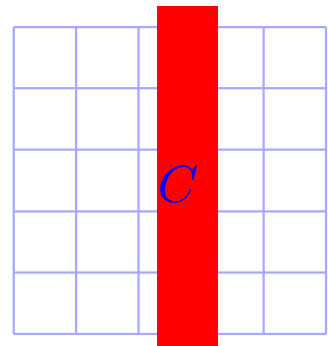
apply  $L$  CNOTs  
in parallel



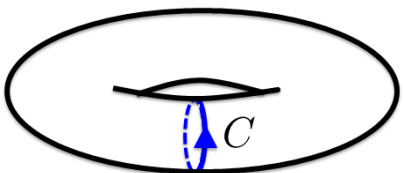
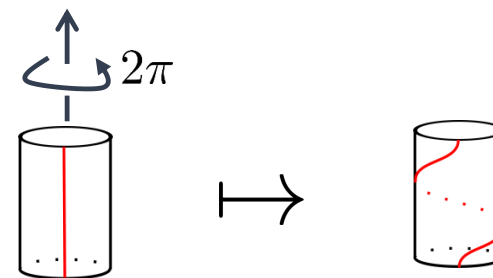
relocate qubits



repeat  $L$  times



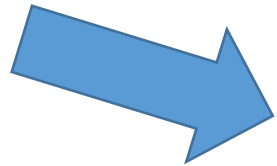
$\Rightarrow$  For every closed, no-contractible loop  $C$ , there is a logical gate  $U(C)$  implementable in depth  $L$



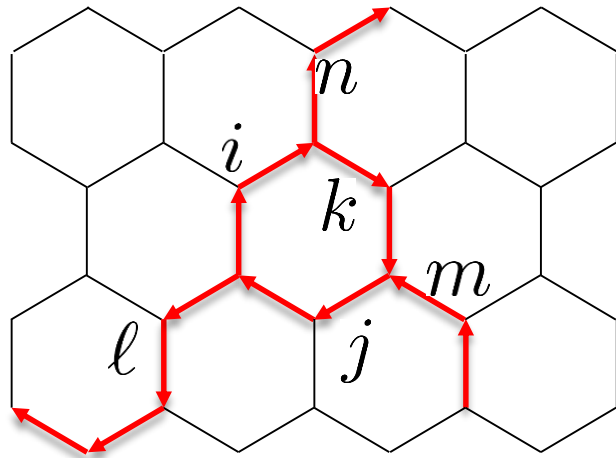
Each  $C$  defines an element  $\vartheta_C \in \text{MCG}$  of the mapping class group of the torus (twisting along  $C$ ).  
 $\vartheta_C \mapsto U(C)$  gives a (projective) representation of  $\text{MCG}$

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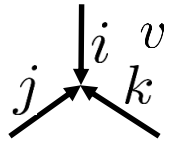
# The Levin-Wen/Turaev-Viro code



local Hilbert space  $\mathbb{C}^d$   
associated to every edge

vertex operator:

$$A_v = \sum_{(i,j,k) \text{ allowed}} |ijk\rangle \langle ijk|$$



plaquette operator:

$$B_p = \frac{1}{\mathcal{D}^2} \sum_{\vec{k}, \vec{k}', \vec{m}} \sum_i d_i \left( \prod_{t=1}^r F_{ik'_{t-1}(k'_t)^*}^{m_t k_t^* k_{t-1}} \right) |\vec{k}', \vec{m}\rangle \langle \vec{k}, \vec{m}|$$

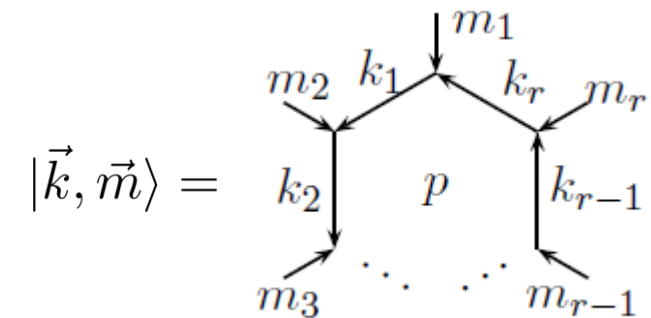
Code space

$$\mathcal{L} \subset (\mathbb{C}^d)^{\otimes N}$$

$$\mathcal{L} = \{ |\Psi\rangle \mid B_p |\Psi\rangle = |\Psi\rangle \forall p, A_v |\Psi\rangle = |\Psi\rangle \forall v \}$$

ingredients:

- finite set of “particle labels”
- involution operation on particle labels
- set of allowed triples
- scalars and a tensor





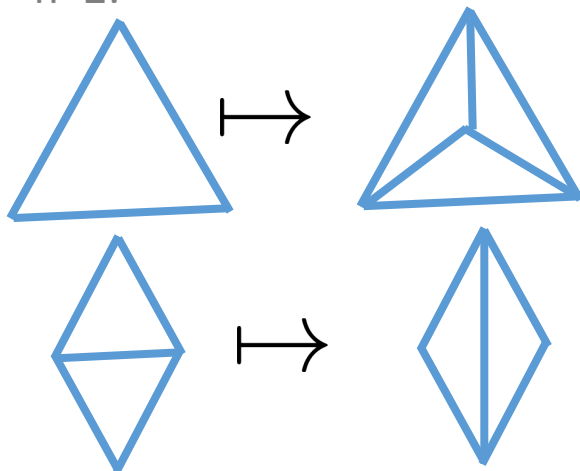
# Manifold-invariants from triangulations

Consider **closed n-manifolds modulo homeomorphism**

**FACT:** For  $n=2,3$ , every equivalence class has a triangulated representative.

**FACT (Pachner):**  $n$ -manifolds homeomorphic  $\iff$   
triangulations related sequence of Pachner moves.

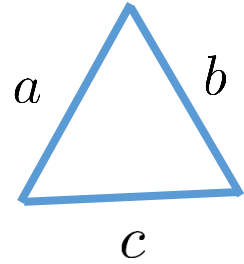
Pachner moves: finite  
list of local changes of  
triangulation, e.g., in  
 $n=2$ :



**Recipe for constructing invariants:**

- associate scalar to every triangulation
- show invariance under Pachner moves

# Example: State-sum invariants



$$\mapsto F_{abc}$$

associate scalar with  
(colored) triangle

define invariant by summing over edge colorings:

$$I(M) = \mathcal{D}^{-\#\text{triangles}} \sum_{\phi} \prod_{\text{triangles } t} g_t^{\phi}$$

triangulated  
2-manifold

sum over all  
colorings

Compatibility with Pachner moves

$$I(\triangle) = I(\text{triangulated triangle})$$

$$I(\text{diamond}) = I(\text{diamond})$$

is equivalent to algebraic conditions

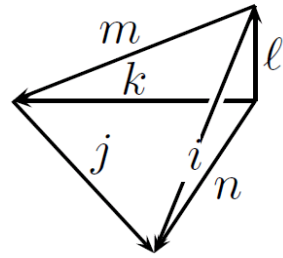
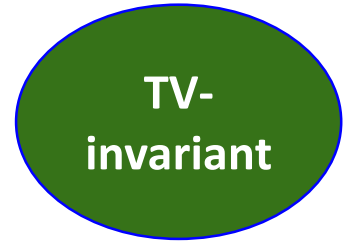
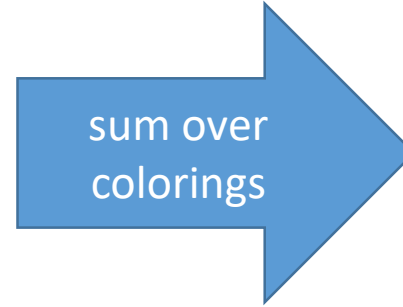
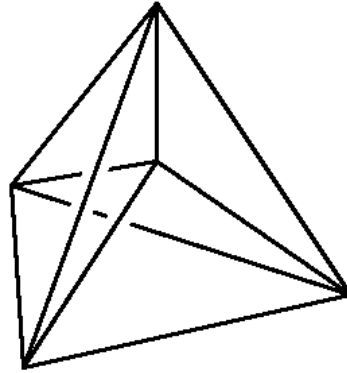
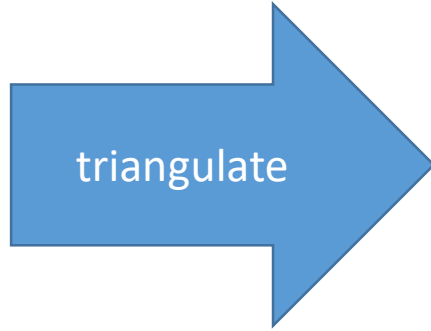
$$\mathcal{D}^{-1} F_{abc} = \mathcal{D}^{-3} \sum_{x,y,z} F_{axz} F_{xby} F_{zyc}$$

$$\sum_x F_{abx} F_{cxd} = \sum_y F_{ayc} F_{dyb}$$

# The Turaev-Viro 3-manifold invariant



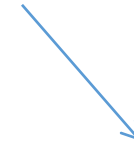
**3-manifold**  
(closed)



$\mapsto$

$$\frac{F_{kl^*n}^{i^*jm}}{\sqrt{d_m d_n}}$$

scalar associated with  
(colored) tetrahedron

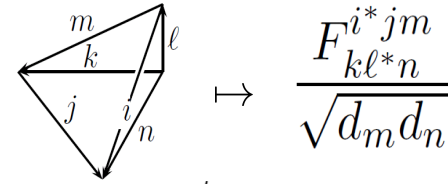


$$\text{TV}(M) = \mathcal{D}^{-2|V_M|} \sum_{\text{colorings } \phi} \prod_{\text{edges } e} d_{\phi(e)} \prod_{\text{tetrahedra } t} g_t^\phi$$

sum over all "allowed" colorings



# Algebraic conditions for invariance (via Pachner moves)



$$\mathrm{TV}_{\mathcal{C}}(M) = \mathcal{D}^{-2|V_M|} \sum_{\text{colorings } \phi} \prod_{\text{edges } e} d_{\phi(e)} \prod_{\text{tetrahedrat}} g_t^{\phi}$$

If

$$1 = 1^*$$

$$d_1 = 1$$

$$d_i = d_{i^*}$$

$$\mathcal{D} = \sqrt{\sum_i d_i^2}$$

$$d_i d_j = \sum_k \delta_{ijk} d_k$$

$$\sum_m \delta_{ijm^*} \delta_{mkl^*} = \sum_m \delta_{jkm^*} \delta_{iml^*}$$

$$F_{kln}^{ijm} \delta_{ijm} \delta_{klm^*} = F_{kln}^{ijm} \delta_{iln} \delta_{jkn^*}$$

$$\sum_n F_{kpn}^{mlq} F_{mns}^{jip^*} F_{lkr}^{jns} = F_{q^*kr}^{jip^*} F_{mls}^{r^*iq^*}$$

$$(F_{kln}^{ijm})^* = F_{k^*l^*n^*}^{i^*j^*m^*}$$

$$F_{kln}^{ijm} = F_{lkn^*}^{jim} = F_{jin}^{lkm^*} = F_{k^*nl}^{imj} \sqrt{\frac{d_m d_n}{d_j d_l}}$$

$$F_{j^*jk}^{ii^*1} = \sqrt{\frac{d_k}{d_i d_j}} \delta_{ijk}$$

\*: involution on set of colors

1: special color

$$\delta_{ijk} \in \mathbb{N} \cup \{0\}$$

$$F_{kln}^{ijm} \in \mathbb{R}$$

$$d_i \in \mathbb{R}_{>0}$$

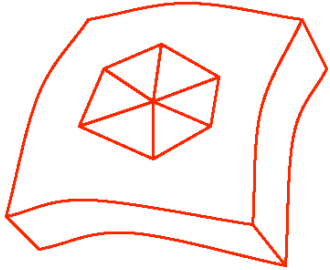
then  $\mathrm{TV}_{\mathcal{C}}$  is a 3-manifold invariant

A spherical category  $\mathcal{C}$  is/provides a solution to these equations.

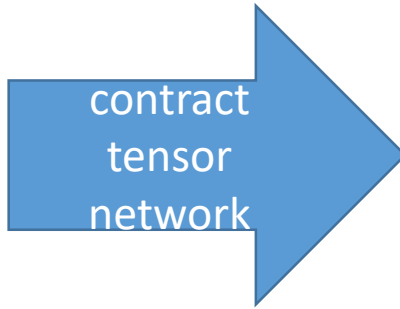
(Barrett and Westbury, hep-th/9311155)

# The Turaev-Viro code $\subset (\mathbb{C}^d)^{\otimes |E|} \cong$ edge colorings of surface triangulation

$$\Sigma \times [-1, 1]$$



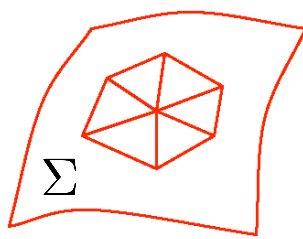
(extend triangulation from  $\Sigma \times \{\pm 1\}$ )



$$(\mathbb{C}^d)^{\otimes |E|}$$



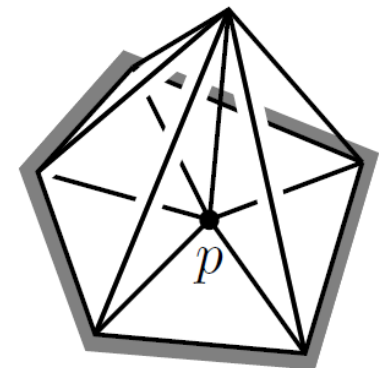
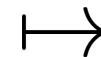
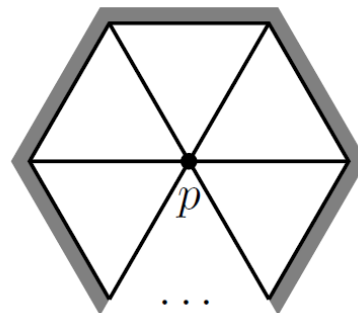
$$(\mathbb{C}^d)^{\otimes |E|}$$



**Turaev-Viro code:** support of this projection in the Hilbert space  $(\mathbb{C}^d)^{\otimes |E|}$

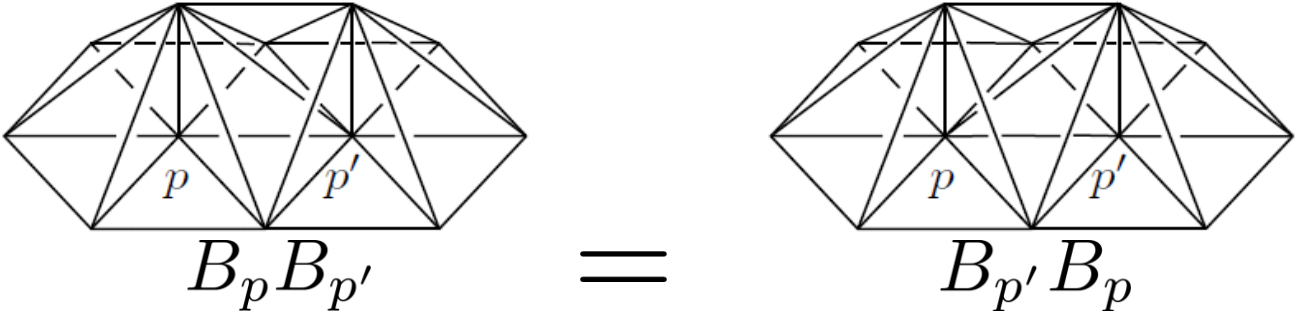
**Local stabilizers:** attaching blisters - set of local operators which are

- projections
- mutually commuting
- stabilize code space

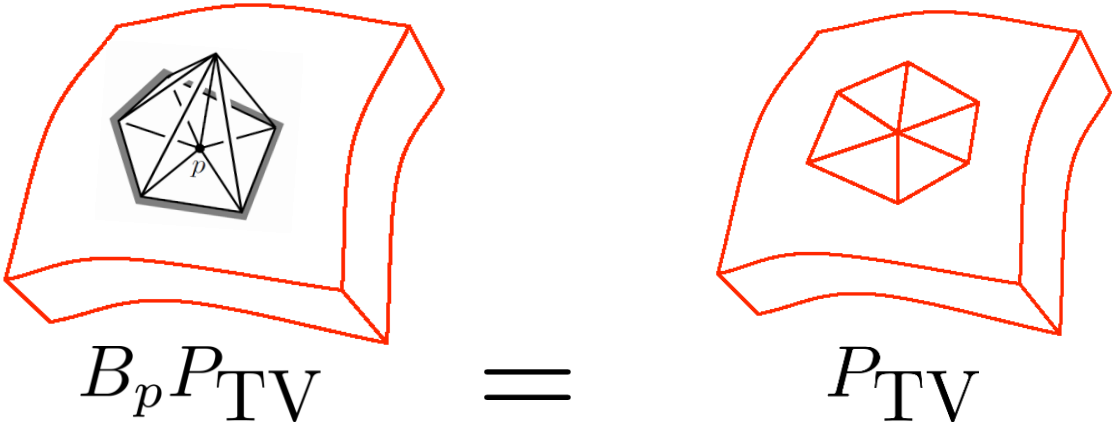


# Blisters: properties from (manifold)invariance

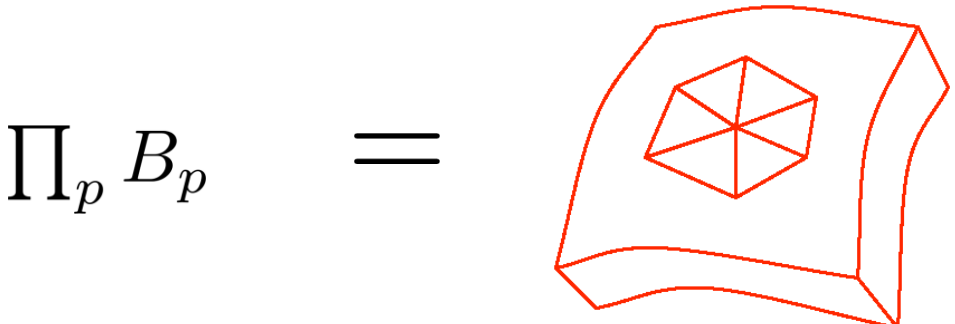
commuting:



stabilize code space:



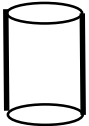
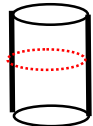
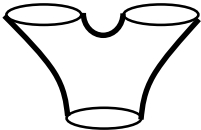
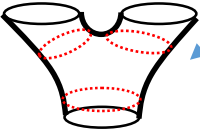
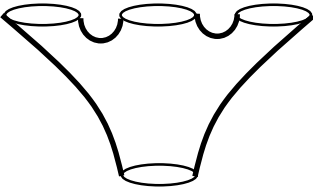
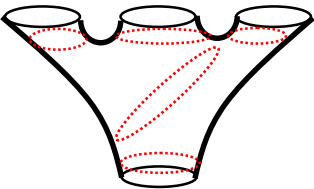
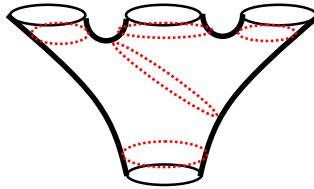
project onto code space



The code space of the Turaev-Viro code

# “Standard bases” from maximal sets of commuting observables

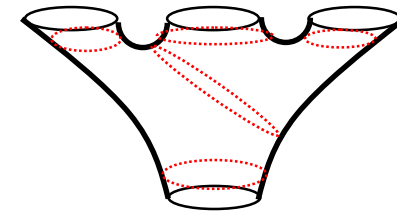
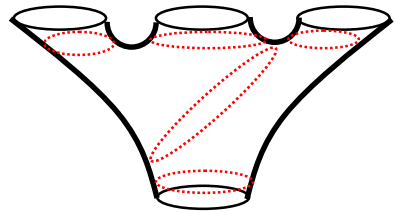
Any DAP-decomposition correspond to a “complete set of observables” and defines a basis of the code space.

surface	DAP-decomposition(s)	elements of standard basis/bases
		$ a\rangle$
		$\begin{array}{c} a \quad b \\ \diagdown \quad / \\ c \end{array}$
		$\begin{array}{c} a_1 \quad a_2 \quad a_3 \\ \diagdown \quad / \quad \diagdown \\ h \quad c \end{array}$
	<p data-bbox="876 968 1447 1011">&lt;- analogy to three spin-1/2s:</p> $(\vec{S}_1 + \vec{S}_2)^2$ $(\vec{S}_1 + \vec{S}_2 + \vec{S}_3)^2$ $S_{\text{total}}^Z$	
		$\begin{array}{c} a_1 \quad a_2 \quad a_3 \\ \diagdown \quad / \quad \diagdown \\ h' \quad c \end{array}$
	$(\vec{S}_2 + \vec{S}_3)^2$ $(\vec{S}_1 + \vec{S}_2 + \vec{S}_3)^2$ $S_{\text{total}}^Z$	

use idempotents of the Verlinde algebra for each loop



F-move: basis change between bases associated with different DAP-decompositions



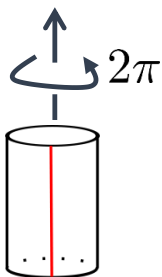
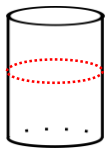
$$\begin{array}{c} a_1 \quad a_2 \quad a_3 \\ \diagdown \quad \diagup \quad \diagdown \\ h \quad \quad \quad c \end{array} = \sum_{h'} F^{a_2 a_1 h}_{c a_3 h'} \begin{array}{c} a_1 \quad a_2 \quad a_3 \\ \diagdown \quad \diagup \quad \diagdown \\ h' \quad \quad \quad c \end{array}$$

some (controlled) unitary  $U(a_2, a_1, a_3, c)_{h, h'}$

....analogous to spin-1/2- 6j symbols

# Mapping class group (generators) and basis elements

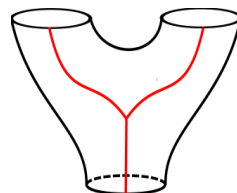
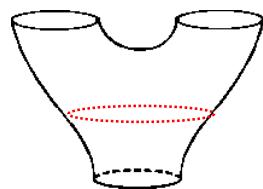
Dehn-twist:



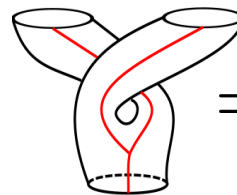
$D$



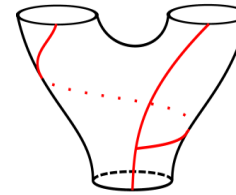
Braid-move:



$B$



$=$



Note: **this** is just a fancy way of writing equation

$$D|i\rangle = \theta_i|i\rangle$$

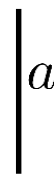
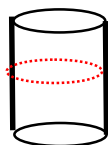
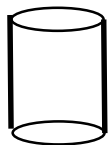
D=twist

surface

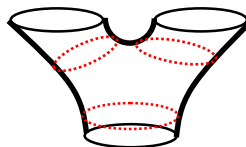
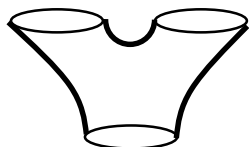
DAP-decomposition(s)

elements of standard basis/bases

**topological phase**



$$|i\rangle = \theta_i |i\rangle$$



**R-matrix**

$$\text{braid} = R_c^{ab} \text{Y-junction}$$

$$B|b, a; c\rangle = R_c^{ab}|a, b; c\rangle$$

B=braid

# Conditions for MCG-representations:

(Moore and Seiberg)

- Consistency of basis changes:

$$\sum_n F_{kpn}^{mlq} F_{mns}^{jip^*} F_{lkr}^{jns} = F_{q^*kr}^{jip^*} F_{mls}^{r^*iq^*}$$

(pentagon-identity)

- Compatibility of basis changes with action of braiding generators:

$$R_m^{ki} F_{lj^*g}^{k^*i^*m} R_g^{kj} = \sum_n F_{lj^*n}^{i^*k^*m} R_\ell^{kn} F_{lk^*g}^{j^*i^*n}$$

$$\theta_i = (R_1^{i^*i})^* \quad (\text{hexagon-identity})$$

- unitarity of representation:

.....

spherical

braided

modular

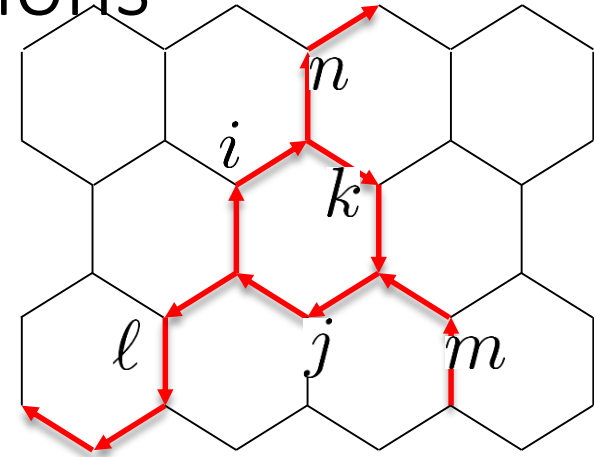
category

Basis states for the Turaev-Viro code

# Levin-Wen ground space and local relations

qudit lattice Hamiltonian

$$H = - \sum_p \text{[blue hexagon with 6 legs]} - \sum_v \text{[green vertex with 3 legs]}$$

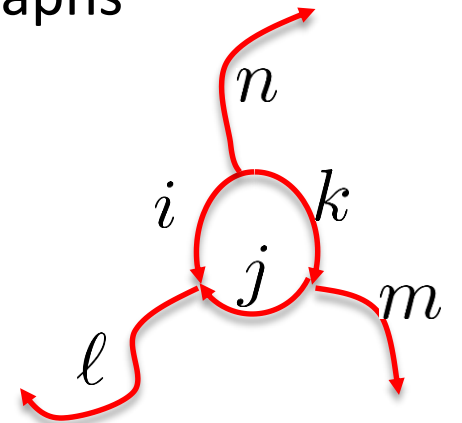


ground state coefficients in computational basis satisfy discrete local “skein” relations, e.g.,

$$\Phi \left( \text{[hexagon with red arrows]} \right) = \Phi \left( \text{[hexagon with red arrows]} \right) \quad \Phi \left( \text{[hexagon with red arrows]} \right) = d_i \Phi \left( \text{[hexagon]} \right)$$

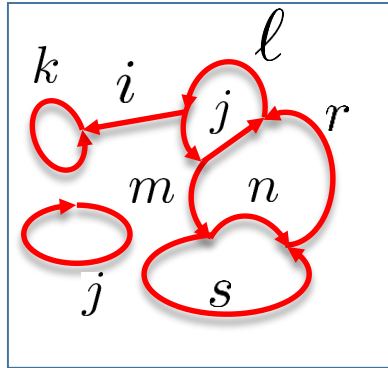
**Consequence:** Ground space is isomorphic to Hilbert space of ribbon graphs (“pictures”) modulo local equivalence relations

ribbon graph space  $\mathcal{H}_\Sigma$



# Ribbon graphs Hilbert space $\mathcal{H}_\Sigma$ for general category

trivalent *labeled directed* graphs (with loops) embedded in  $\Sigma$



**State:** formal linear combination of ribbon graphs

$$\alpha \left[ \text{graph with loops } k, j, i \right] + \beta \left[ \text{graph with loop } k \right] + \gamma \left[ \text{graph with loops } k, j, m, n, r, s \right] + \dots$$

modulo **local relations**

$$\left( \text{loop } i \right) = \left( \text{loop } i \right)$$

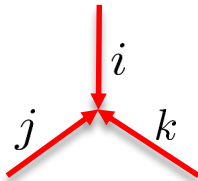
$$\left( \text{loop } i \right) = d_i \quad \text{q-dimensions}$$

$$-i \rightarrow \left( \text{loop } j \right) = 0$$

$$\left( \text{trivalent vertex } i, j, k \text{ with loop } m \right) = \sum_n F_{kln}^{ijm} \left( \text{trivalent vertex } i, j, k \text{ with loop } n \right) \quad \text{F-symbol}$$

**fusion rules**

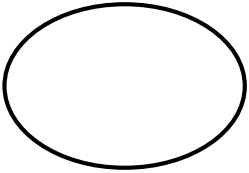
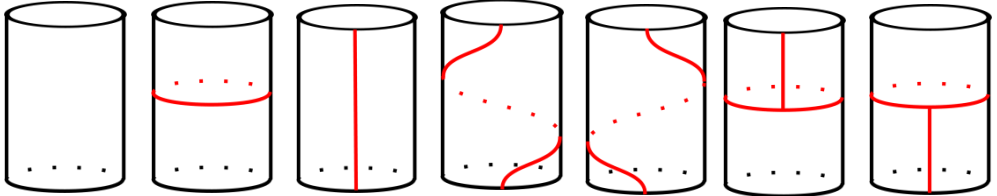
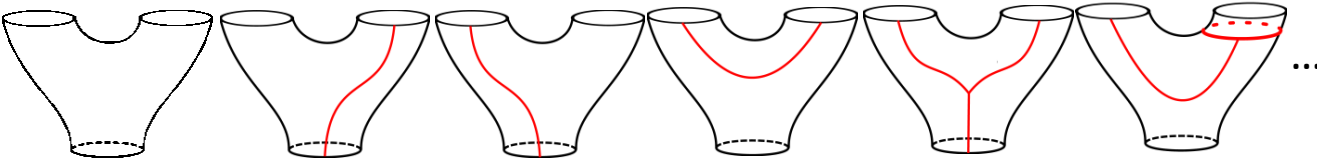
(set of allowed triples):



$$\text{dual labels: } \begin{array}{c} i \\ \longrightarrow \end{array} = \begin{array}{c} \longleftarrow \\ i^* \end{array}$$

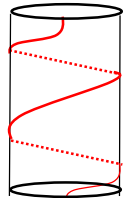
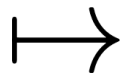
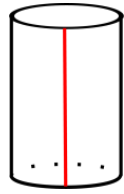
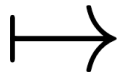
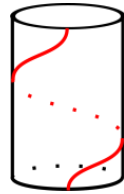
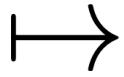
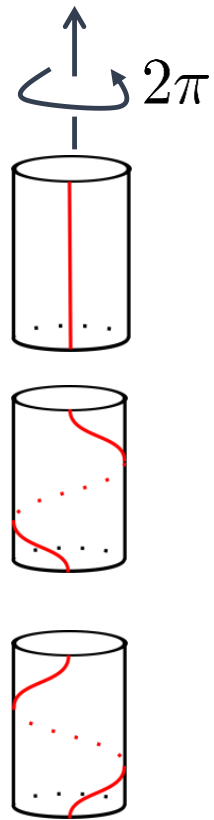
$$\text{trivial label (absence of string): } \begin{array}{c} 1 \\ \longrightarrow \end{array} = \begin{array}{c} \longleftarrow \\ 1 \end{array} =$$

# Ribbon graph bases of $\mathcal{H}_\Sigma$ for *Fib*

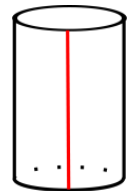
Surface $\Sigma$	$\dim \mathcal{H}_\Sigma$	Example basis
<b>Disc</b> (1-punctured sphere)	<b>1</b>	
<b>Annulus</b> (2-punctured sphere)	<b>7</b>	
<b>Pair of pants</b> (3-punctured sphere)	<b>65</b>	
n-punctured sphere	$2^{\Omega(n)}$	

Next: Description of bases compatible with action of (generators of) mapping class group!

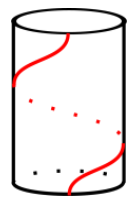
# Action of Dehn twist on $\mathcal{H}_{\Sigma_2}$ for *Fib*



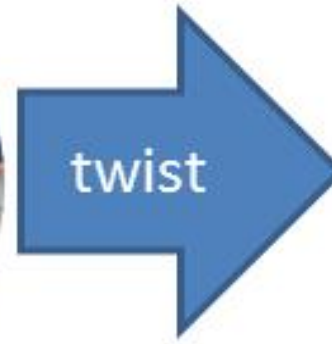
$$= \frac{1}{\phi}$$



$$- \frac{1}{\phi}$$

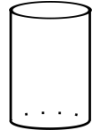
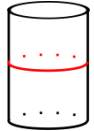
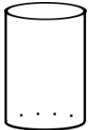
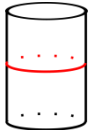

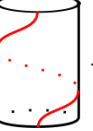
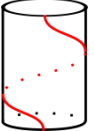

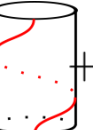
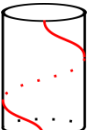

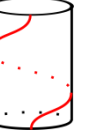

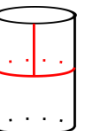
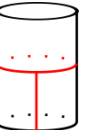


$$+$$



**Goal:** identify “fusion tree basis” (eigenvectors of twist)



Eigenvector	eigenvalue (twist)	name	boundary labels
 + $\phi$ 	1	$1 \otimes 1$	} 1, 1
 + $\frac{1}{\phi}$ 	1	$\tau \otimes \tau$	
 + $\phi e^{-3\pi i/5}$  + $\phi e^{3\pi i/5}$ 	$e^{4\pi i/5}$	$1 \otimes \tau$	} $\tau, \tau$
 + $\phi e^{3\pi i/5}$  + $\phi e^{-3\pi i/5}$ 	$e^{-4\pi i/5}$	$\tau \otimes 1$	
$\phi$  +  + 	1	$\tau \otimes \tau$	} $\tau, 1$
	1	$\tau \otimes \tau$	
	1	$\tau \otimes \tau$	} 1, $\tau$

Anyonic fusion basis obtained by diagonalization

fusion space basis element

$$V_i^i = \mathbb{C} \left| i \right.$$

topological phase

$$\mathcal{J}_i = \theta_i \left| i \right.$$

anyon type

$$i \text{ of "doubled" theory } Fib \otimes Fib'$$

multiplicity index

for different realizations as subspaces of  $\mathcal{H}_{\Sigma_2}$

# Anyonic fusion basis from “doubled” manifold $\Sigma \times [-1, 1]$

**Goal:** find anyonic fusion basis states on  $\Sigma$

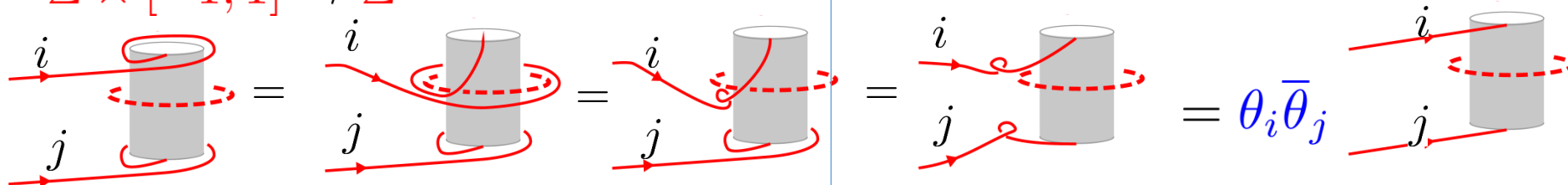
Intermediate step: identify relevant ribbon graphs on  $\Sigma \times [-1, 1]$

A recipe which does not involve diagonalization

simple derivation of topological phase:

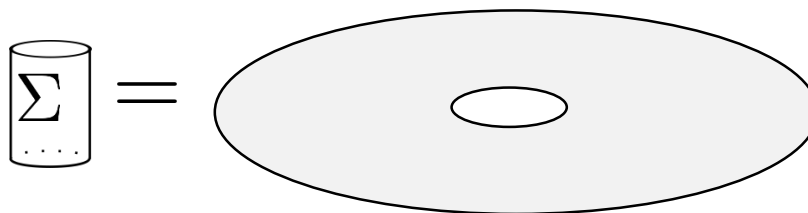
Map ribbon graphs  $\Sigma \times [-1, 1] \rightarrow \Sigma$

using “vacuum” lines

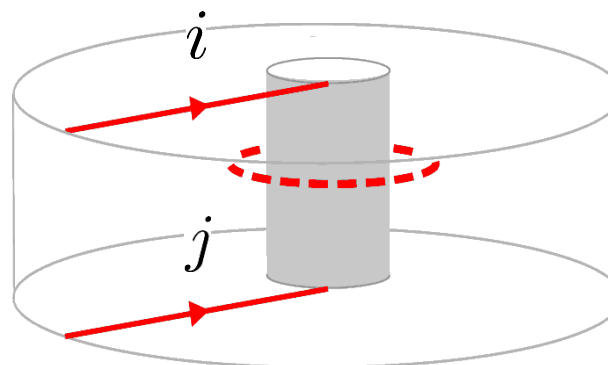


Example: find element for annulus

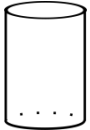

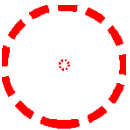

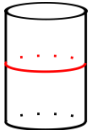

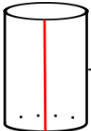




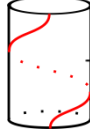
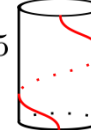

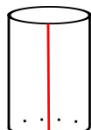
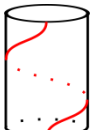


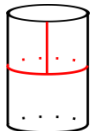

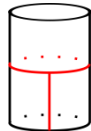

$$\uparrow i \otimes j$$



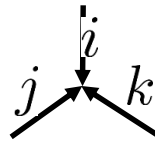
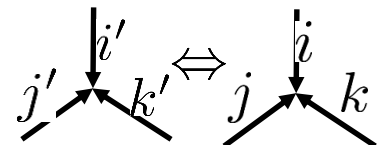
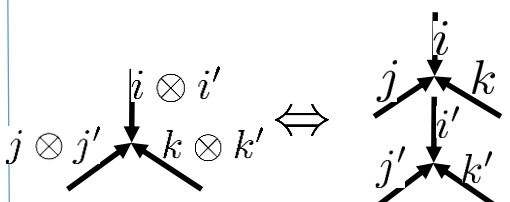
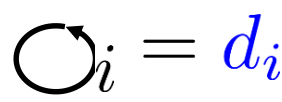
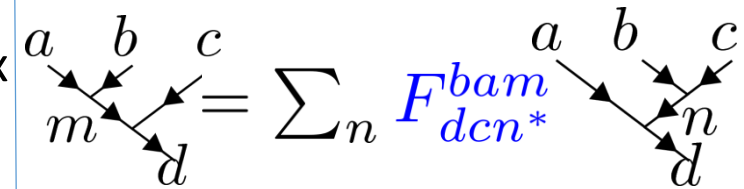
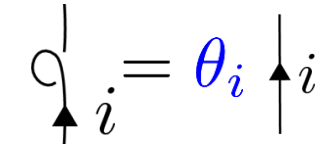
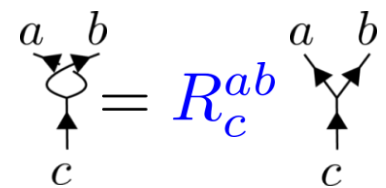
some ribbon graph on



$$| \cdot \rangle := \frac{1}{\sum_i d_i^2} \sum_j d_j | j \rangle$$

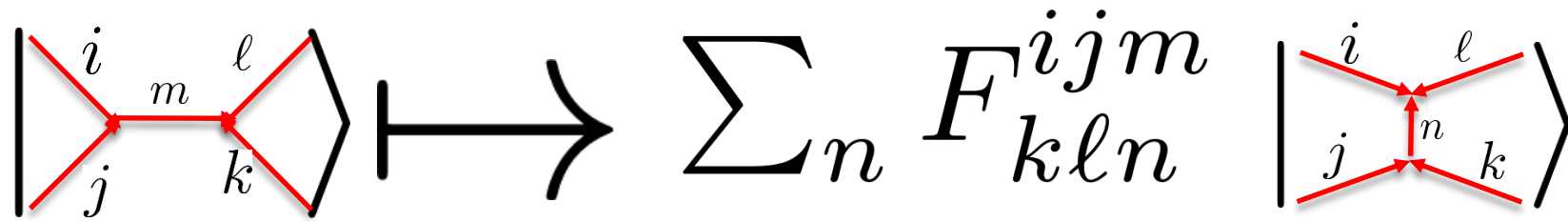
eigenvector		3D-representation	name	boundary labels
	 $+ \phi$ 		$1 \otimes 1$	} 1, 1
	 $+ \frac{1}{\phi}$ 		$\tau \otimes \tau$	
	 $+ \phi e^{-3\pi i/5}$  $+ \phi e^{3\pi i/5}$ 		$1 \otimes \tau$	} $\tau, \tau$
	 $+ \phi e^{3\pi i/5}$  $+ \phi e^{-3\pi i/5}$ 		$\tau \otimes 1$	
$\phi$	 $+$  $+$ 		$\tau \otimes \tau$	
			$\tau \otimes \tau$	} $\tau, 1$
			$\tau \otimes \tau$	

# Derived categories: basic data

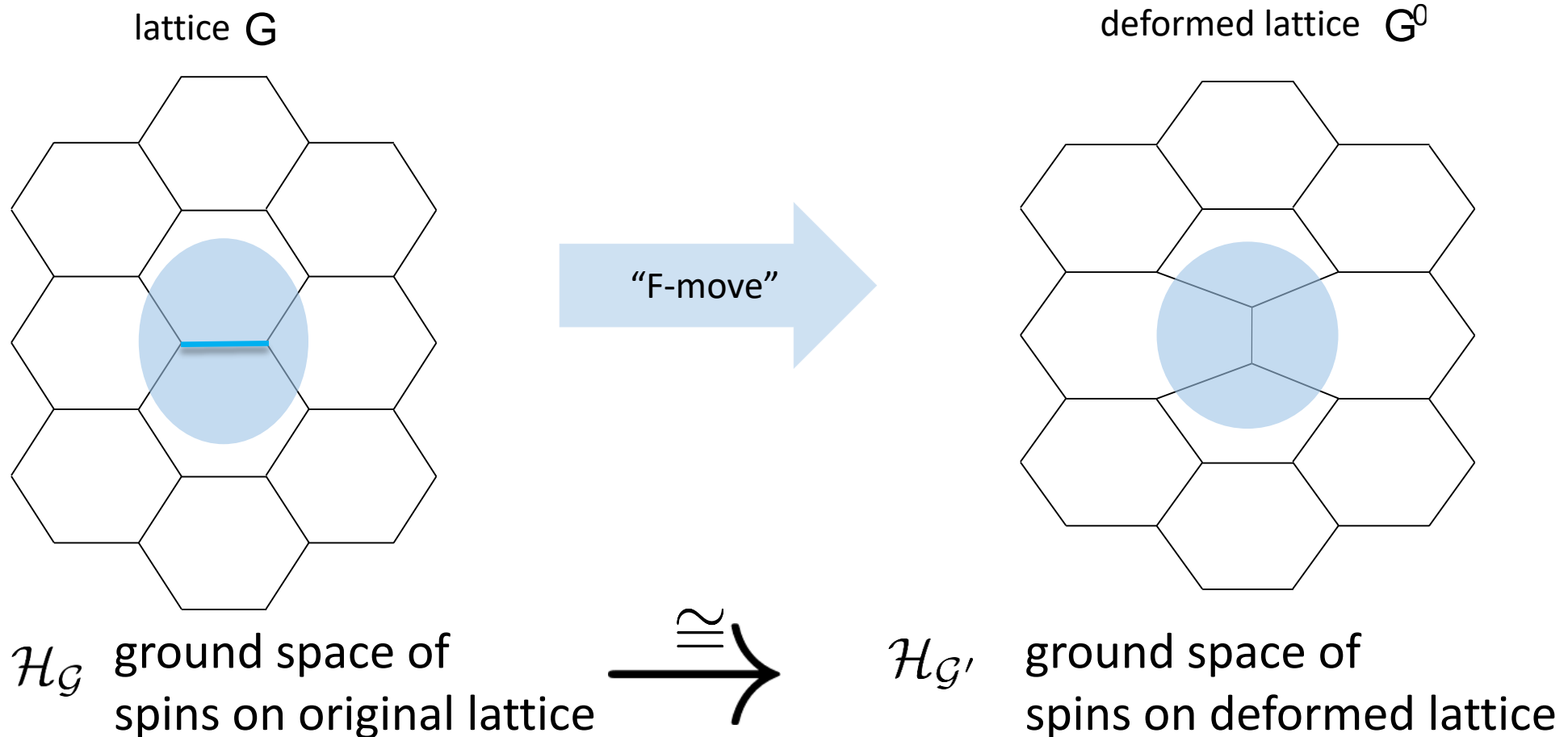
	modular tensor category $\mathcal{C}$	dual category $\mathcal{C}'$	doubled category $\mathcal{C} \otimes \mathcal{C}'$
Unitary, braided, semisimple, *			
Particles	$\{1, i, j, \dots\}, *$	$\{i' \mid i \in \mathcal{C}\}$	$\{i \otimes j' \mid \begin{matrix} i \in \mathcal{C}, \\ j' \in \mathcal{C}' \end{matrix}\}$
Fusion rules	 (set of) allowed triples		
q-dim	 $= d_i$	$d_{i'} = d_i$	$d_{i \otimes j'} = d_i d_{j'}$
F-matrix	 $= \sum_n F_{dcn}^{bam}$	$F_{d'c'n'}^{b'a'm'} = F_{dcn}^{bam}$	$F \otimes F'$
top. phase	 $= \theta_i$	$\theta_{i'} = \bar{\theta}_i$	$\theta_{i \otimes j'} = \theta_i \theta_{j'}$
R-matrix	 $= R_c^{ab}$	$R_{c'}^{a'b'} = \overline{R_c^{ab}}$	$R \otimes R'$

# Computation with Turaev-Viro codes

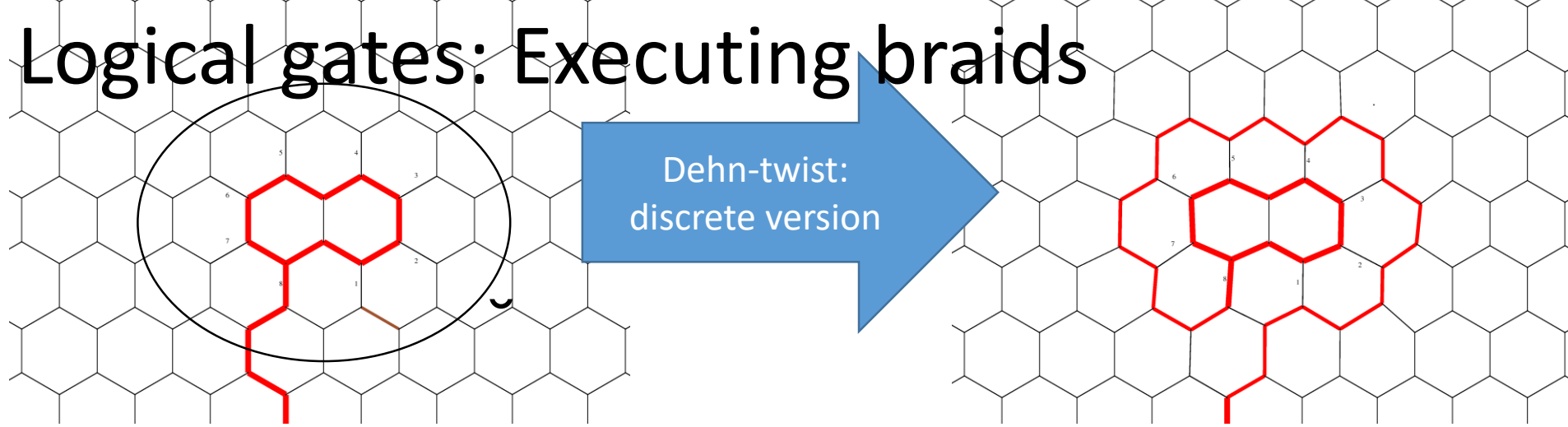
# Different lattices and F-move isomorphism



For unitary tensor categories, this is a **unitary 5-qudit gate**.

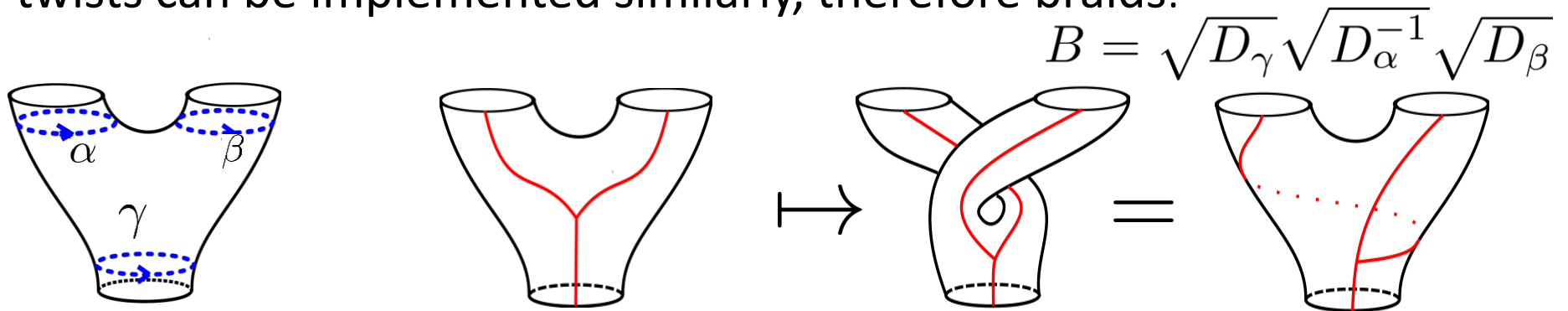


# Logical gates: Executing braids



Can be implemented by sequence of  $O(|\gamma|^2)$  F-moves (5-qudit gates)

$\pi$ -twists can be implemented similarly, therefore braids:



## universal gate set:

- braids generate dense subgroup of unitaries on subspace of  $\mathcal{H}_\Sigma$  for (doubled) Fib
- for appropriate encoding, approximation of universal gate set by Solovay-Kitaev (Freedman, Larsen, Wang'02)

# Gate sets obtained from the mapping class group



<i>TQFT</i>	<i>mapping class group (braiding)</i> <i>contained in</i>
$D(\mathbb{Z}_2)$	Pauli group
abelian anyon model	generalized Pauli group
Fibonacci model	universal
Ising model	Clifford group
generic anyon model	model-dependent
generic anyon model	universal



# Conclusions and open problems

- Turaev-Viro codes offer a rich class of examples for potential platforms for topological quantum computation.
- The mapping class group representation can be “decomposed” using the string-net formalism
- Explicit constructions of protected/transversal gates for TQFTs?
- Performing syndrome-measurement & error correction, thresholds for fault-tolerance?
- Higher-dimensional generalizations?