

A collection of results on polynomial maps over finite fields

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Basics

Let R be a ring. Denote:

- $MA_n(R)$ the set of polynomial endomorphisms,
- $GA_n(R)$ the set of polynomial automorphisms,
- $BA_n^0(R)$ is the set of strictly upper triangular polynomial automorphisms,
- $TA_n(R) := \langle BA_n^0(R), GL_n(R) \rangle$ the set of tame polynomial automorphisms,
- $SA_n(R) = \{F \in GA_n(R) \mid \det(\text{Jac}(F)) = 1\}$,
- $STA_n(R) = TA_n(R) \cap SA_n(R)$.

Let $q = p^m$ where p is prime. We can define

$$\pi_q : MA_n(\mathbb{F}_q) \longrightarrow \text{Maps}((\mathbb{F}_q)^n, (\mathbb{F}_q)^n)$$

and thus also

$$\pi_q : GA_n(\mathbb{F}_q) \longrightarrow \text{Perm}((\mathbb{F}_q)^n).$$

Main question

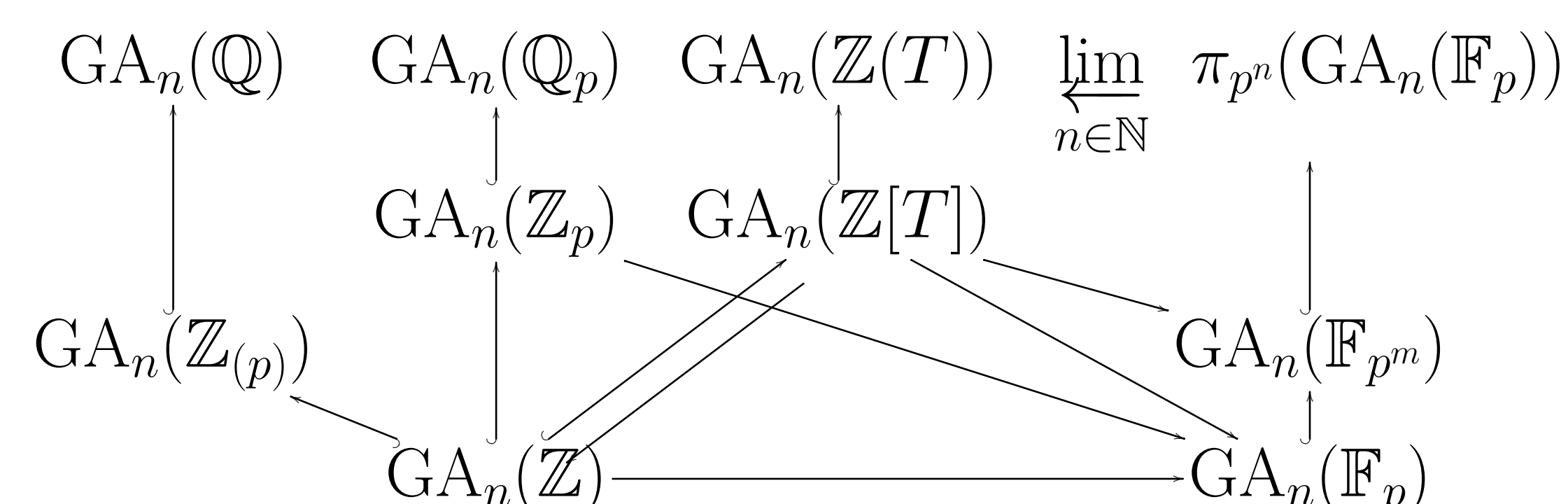
What is $\varprojlim_{m \in \mathbb{N}} (GA_n(\mathbb{F}_q), \pi_{q^m}(TA_n(\mathbb{F}_q)))$ and are they different?

Finding a difference would imply that there exist wild polynomial automorphisms.

Theorems on the case $m = 1$

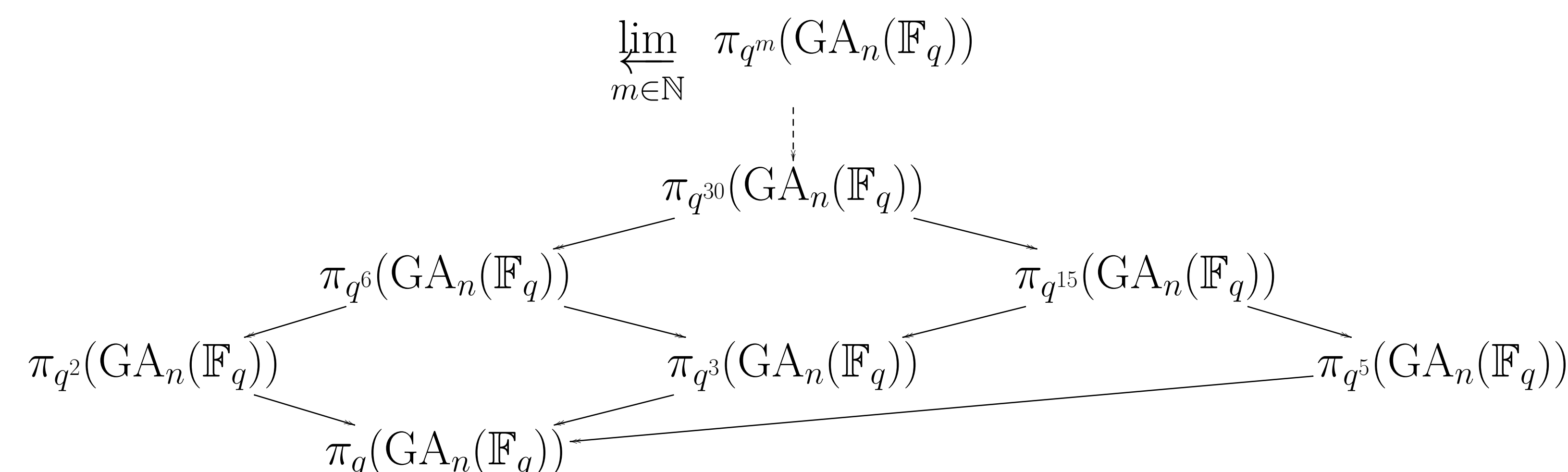
- $\pi_q TA_n(\mathbb{F}_q) = \text{Sym}((\mathbb{F}_q)^n)$ if $q = \text{odd}$ or $q = 2$, and
- $\pi_q TA_n(\mathbb{F}_q) = \text{Alt}((\mathbb{F}_q)^n)$ if $q = \text{even}$ but not $q = 2$.
- $\pi_q STA_n(\mathbb{F}_q) = \text{Alt}((\mathbb{F}_q)^n)$,
- unless $q = 2$, when it is $\text{Sym}((\mathbb{F}_q)^n)$.

Interesting connections



The profinite polynomial automorphism group

Since there exist restriction maps $\pi_{q^m} GA_n(\mathbb{F}_q) \longrightarrow \pi_q GA_n(\mathbb{F}_q)$ we get the following chain and inverse limit:



We call $\varprojlim_{m \in \mathbb{N}} \pi_{q^m}(GA_n(\mathbb{F}_q))$ the profinite polynomial automorphism group (which contains $GA_n(\mathbb{F}_q)$). Similarly, we define the profinite tame automorphism group $\varprojlim_{m \in \mathbb{N}} \pi_{q^m}(TA_n(\mathbb{F}_q))$ and profinite polynomial endomorphisms

$$\varprojlim_{m \in \mathbb{N}} \pi_{q^m}(MA_n(\mathbb{F}_q)).$$

Theorem: Wild automorphisms in profinite tame group

Assume

- (1) $F \in GA_n(\mathbb{F}_q[X_{n+1}])$, (2) $F \in TA_n(\mathbb{F}_q(X_{n+1}))$, (3) $F(X_{n+1} = c) \in TA_n(\mathbb{F}_q)$ for all $c \in \mathbb{F}_q$.

Then F is in the profinite tame automorphism group, i.e.

$$F \in \varprojlim_{m \in \mathbb{N}} \pi_{q^m}(TA_n(\mathbb{F}_q)).$$

In particular:

$$GA_2(\mathbb{F}_q[Z]) \subseteq \varprojlim_{m \in \mathbb{N}} \pi_{q^m}(TA_n(\mathbb{F}_q)).$$

This theorem implies that it is not possible to distinguish for example Nagata's automorphism from a tame automorphism by only examining its permutations.

A theorem on the Derksen group

If $n \geq 3$, define $DA_n(\mathbb{F}_q) = \langle \text{Aff}_n(\mathbb{F}_q), E \rangle$ where

$$E = (x_1 + (x_1 x_3 \cdots x_n)^{p-1}, x_2, \dots, x_n).$$

This group we called the Derksen group. Theorem:

$$\varprojlim_{m \in \mathbb{N}} \pi_{q^m}(DA_n(\mathbb{F}_q)) = \varprojlim_{m \in \mathbb{N}} \pi_{q^m}(TA_n(\mathbb{F}_q))$$

so we do have actual smaller groups that give the same profinite groups. Well - as soon as we prove that $DA_n(\mathbb{F}_q)$ is not equal to $TA_n(\mathbb{F}_q)$!

The profinite polynomial endomorphism monoid

We define $\varprojlim_{m \in \mathbb{N}} \pi_{q^m}(MA_n(\mathbb{F}_q))$ as the profinite polynomial endomorphism monoid. Consider

$$\mathcal{M}_{n,m}(\mathbb{F}_q) := \pi_{q^m} MA_n(\mathbb{F}_q) \cap \text{Perm}((\mathbb{F}_q^m)^n).$$

Then $\varprojlim_{m \in \mathbb{N}} \mathcal{M}_{n,m}(\mathbb{F}_q)$ is the subset of invertible elements in $\varprojlim_{m \in \mathbb{N}} \pi_{q^m}(MA_n(\mathbb{F}_q))$, i.e. we can call it the profinite polynomial endomorphism group. How does it look like? Define X as the set of orbits of \mathbb{F}_q^n under the action of $\text{Gal}(\mathbb{F}_q^m : \mathbb{F}_q)$, and let X_d be the set of orbits of size d . Then

$$\varprojlim_{m \in \mathbb{N}} \mathcal{M}_{n,m}(\mathbb{F}_q) \cong \prod_{d \in \mathbb{N}} ((\mathbb{Z}/d\mathbb{Z}) \text{ wr}_{X_d} \text{Perm}(X_d)).$$

Profinite tame group vs. profinite polynomial endomorphism group

How much does $\varprojlim_{m \in \mathbb{N}} \pi_{q^m}(GA_n(\mathbb{F}_q))$ differ from

$\varprojlim_{m \in \mathbb{N}} \mathcal{M}_{n,m}(\mathbb{F}_q)$? By far it is not equal - but: define

$$\Pi_q : GA_n(\mathbb{F}_q) \longrightarrow \text{Perm}(X)$$

then consider $\Pi_{q^m}(TA_n(\mathbb{F}_q))$. Apparently:

$\Pi_{q^m}(TA_n(\mathbb{F}_q)) = \mathcal{M}_{n,m}(\mathbb{F}_q)$ if $n \geq 3$ except finitely many q . In particular: $\Pi_{q^m}(GA_n(\mathbb{F}_q)) = \mathcal{M}_{n,m}(\mathbb{F}_q)$ in those cases!

This gives a foothold in tacking (parts of) the main question!

Alternative to LFIHderivations: \mathbb{Z} -flows

If k a field, then k -actions on k^n correspond to locally nilpotent derivations (LNDs) on k^n if $\text{char } k = 0$. If $\text{char}(k) = p$, then k -actions on k^n correspond to so-called *locally finite iterative higher derivations*. Longer name, less nice properties! For example:

$$(x + y + z, y + z, z)$$

is a unipotent map, but is not exponent of a LFIHD if $\text{char}(k) = 2$ (for $\exp(D)$ has order p). Bah!

Example of a \mathbb{Z} -flow

Define

$$R := \mathbb{Z}[Q_i \mid i \in \mathbb{N}] / (p, Q_i^p - Q_i \mid i \in \mathbb{N})$$

where Q_i corresponds to $\mathbb{Z} \longrightarrow \mathbb{F}_p$ given by $t \longrightarrow \binom{t}{p^i} \text{ mod } p$. Then $F := (x + y + z, y + z, z) \in \text{TA}_3(\mathbb{F}_2)$ has a " \mathbb{Z} -flow":

$$F_t := (x + Q_0 y + (Q_1 + Q_0)z, y + Q_0 z, z).$$

Indeed, $F_t(t = n) = F^n$ for each $n \in \mathbb{Z}$.

Interesting object

This opens up the idea to examine $GA_n(R)$.

Fast forward functions from cryptography

It is desirable of a function f if $f^n(v)$ is efficiently computable w.r.t. computation of $f(v)$ for any n, v . Let $\sigma \in \pi_p(BA_n^0(\mathbb{F}_p))$ such that σ has only one orbit in \mathbb{F}_p^n . Then there exists $\tau \in BA_n^0(\mathbb{F}_p)$, D a diagonal linear map, and a trivial map $\zeta : (\mathbb{F}_p)^n \longrightarrow \mathbb{Z}/p^n\mathbb{Z}$ such that

$$\zeta D \tau \sigma \tau^{-1} D^{-1} \zeta^{-1} = \text{inc}$$

where $\text{inc}(z) = z + 1$ on $\mathbb{Z}/p^n\mathbb{Z}$, making iterations of σ efficiently computable.

References

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