

Carsten Thomassen :

Decompositions and
Grötzsch's theorem

Institutional lecture at

ADONET-CIRM School of
Graphs and Algorithms

Oct. 24, 2007

Steiner triple system
on a set S

is a collection of
triples such that each
pair in S is in
precisely one triple

Theorem A STS
on S exists iff

n is odd and

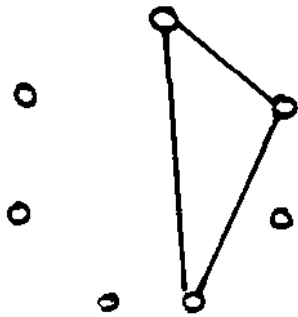
$$\binom{n}{2} \equiv 0 \pmod{3}$$

Set system generalization:

Theory of block design

A STS on the set S is also a decomposition of $K_{|S|}$ into edge-disjoint triangles.

Example: A decomposition of K_7 into triangles



Rotate this triangle to get 7 edge-disjoint triangles.

Graph decomposition

H fixed graph
 G arbitrary graph, for example a big complete graph.

An H -decomposition of G is a collection H_1, H_2, \dots, H_m of pairwise edge-disjoint subgraphs in G such that

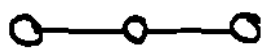
$$H_1 \cup H_2 \cup \dots \cup H_m = G$$

and each

$$H_i \cong H$$

Simplest example:

$H =$ the 2-path P_3



A connected graph G
has a P_3 -decomposition

iff $|E(G)|$ is even.

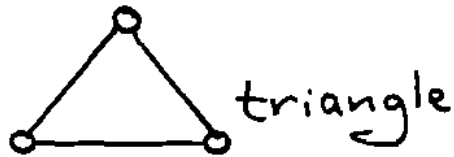
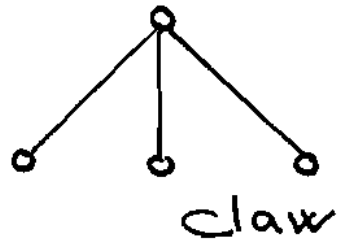
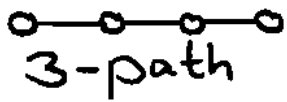
Proof: Delete from
 G the edges of a
2-path such that
the resulting graph
is connected
(or conn, + isolated
vertex)

Dor and Tarsi (1997)

For each fixed
connected graph with
at least 3 edges,
the H -decomposition
problem is
NP-complete.

There are several
results when H is
fixed and
 G is either very dense
or of a very special
structure

Smallest interesting cases



J.-C. Bermond +
Schönheim (1977)

H graph with ≤ 4 vertices

Then the complete graph K_n has an H -decomposition if and only if n satisfies ...

This extends the theorem on Steiner triples.

R. M. Wilson (1975)

H any fixed graph
 n sufficiently large
natural no.

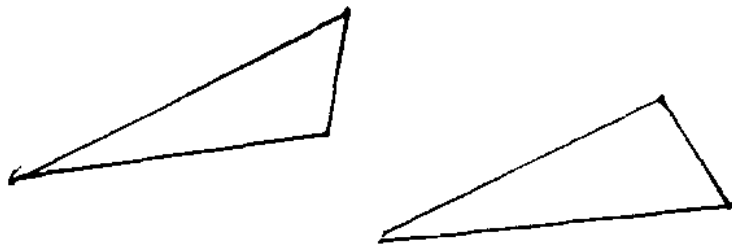
Then K_n has an H -decomposition if and only if

K_n has the right size
($\binom{n}{2}$ divisible by $|E(H)|$)
and the right degrees.

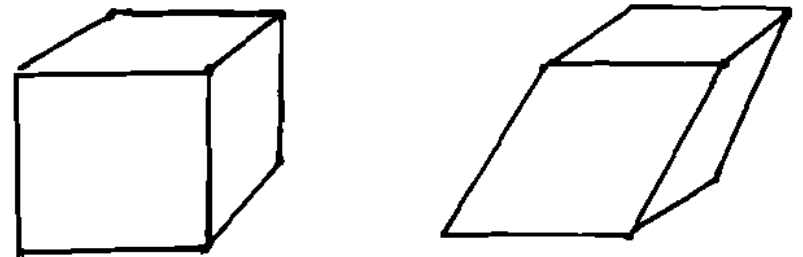
If you fix the sides of a triangle, the triangle is rigid.

More generally :

Two triangles with the same sides are congruent.



If you fix the edges of a polyhedron in \mathbb{R}^3 it need not be rigid.



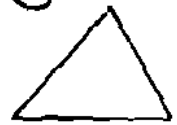
But, if you fix the sides (facets) it is rigid.

Cauchy If two polytopes have congruent facets, they are congruent.

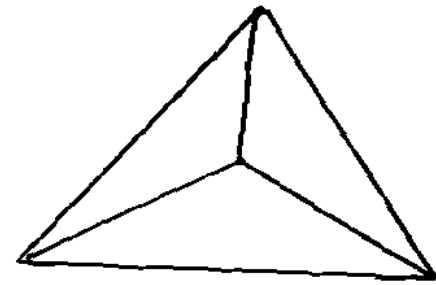
Motivated by Cauchy's
result and
rigidity of convex polyhedra

M. Dehn (1916) proved:

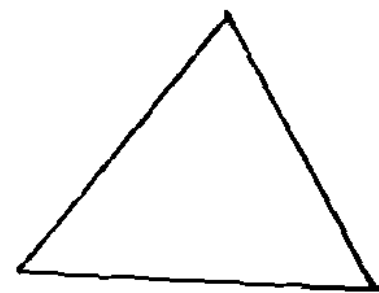
The graph of a
triangulated polyhedron
that is,
a planar triangulation
has an edge-decomposition
into a



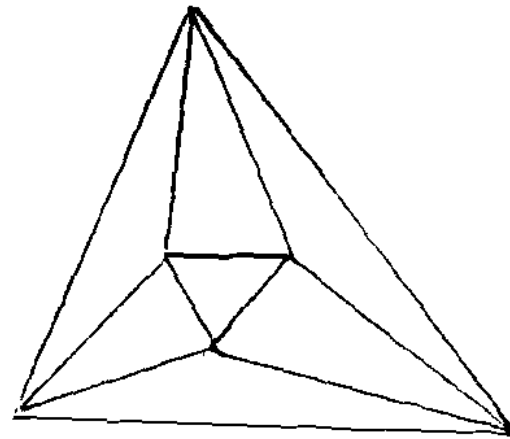
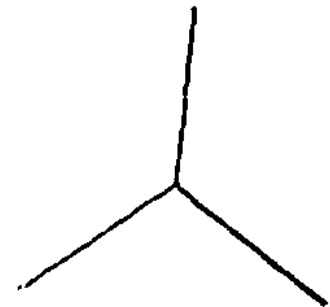
and a collection of
claws:



=



+



A proof of Dehn's result

G graph

$$f: V(G) \rightarrow \{0, 1, 2, \dots\}.$$

Then G has an orientation such that

$$d^+(v) \leq f(v) \quad \forall v \in V(G)$$

iff, for every subgraph H of G ,

$$|E(H)| \leq \sum_{v \in V(H)} f(v).$$

First proved by Halimi (1965).

Follows also from Edmonds' matroid partition theorem

Application to a planar graph G .

Put $f(v_1) = f(v_2) = 0$
and $f(v) = 3$ for all other v .

Then G has an orientation such that

$$d^+(v_1) = d^+(v_2) = 0$$

and $d^+(v) \leq 3$ for all other v .

If G is a triangulation equality holds because

$$|E(G)| = 3n - 6$$

which gives a claw-decomposition.

If G triangulates
the sphere + g handles
then $|E(G)| = 3n - 6 + 6g$

Define $f: V(G) \rightarrow \mathbb{N}$
such that

$f(v) = 6$ for $2g - 2$
vertices and

$f(v) = 3$ for all others.

If $d^+(v) \leq f(v) \forall v$
then equality holds.

Can the $2g - 2$ vertices
be chosen such that
the f -inequality
holds?

We don't know, but:

by another method
we prove:

Barat + CT:

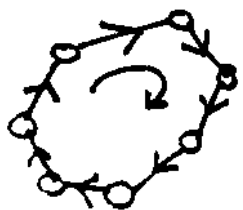
Every triangulation
of every surface has
a claw decomposition

Observation of Jaeger:

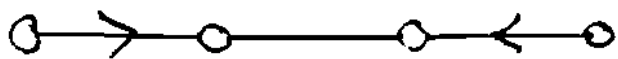
Every cubic 2-conn.
graph has a
3-path-decomposition

Proof: Let M be a
perfect matching.

Then $G-M$ is a collection
of cycles. Make each to
an oriented cycle:



For each edge
in M , take



We have decomposition
results for
dense graphs (Wilson)
and
graphs with special
structure
(Dehn, Jaeger).

Conjecture for general
graphs (Barat + CT)

For every tree T ,
there exists a number
 k_T such that each
 k_T -edge-connected graph
of size divisible by
 $|E(T)|$ has a
 T -decomposition

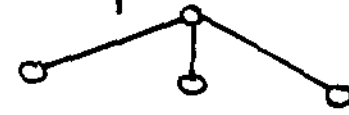
Important that T is a tree because there are 10^{10} - edge - conn. graphs of girth $> 10^{10}$.

When we made the conj. we could not prove it for one single tree of size ≥ 3 .

Recently I proved it for one (but only one) tree for graphs of edge - conn.

10^{10} 10^{14}

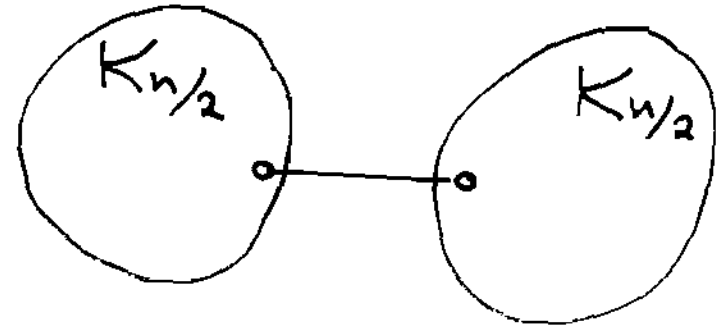
Now focus on the claw



Theorem (Barat + CT)

$$\left. \begin{array}{l} |V(G)| = n \\ \delta(G) > \frac{n}{2} \\ n \text{ large} \\ |E(G)| \equiv 0 \pmod{3} \end{array} \right\} \implies G \text{ has a claw-decompos.}$$

$n/2$ best possible

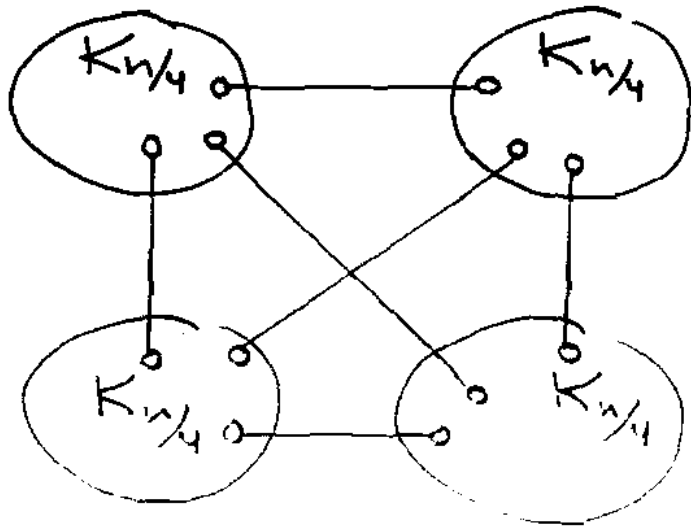


Theorem (JB+CT):

$|V(G)| = n$
 $|E(G)| \equiv 0 \pmod{3}$
 G 2-edge-conn.
 $\delta(G) > \frac{n}{4}$
 n large

$\Rightarrow G$ has a claw-decompos.

$n/4$ best possible



Conjecture (JB+CT)

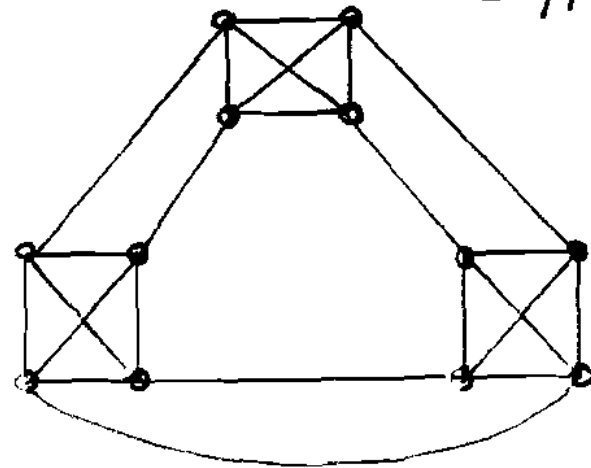
G 10-edge-conn.
 $|E(G)| \equiv 0 \pmod{3}$

$\Rightarrow G$ has a claw-decompos.

What about 9-edge-conn.?
 8
 7
 6
 5

?
 ?
 ?
 ?
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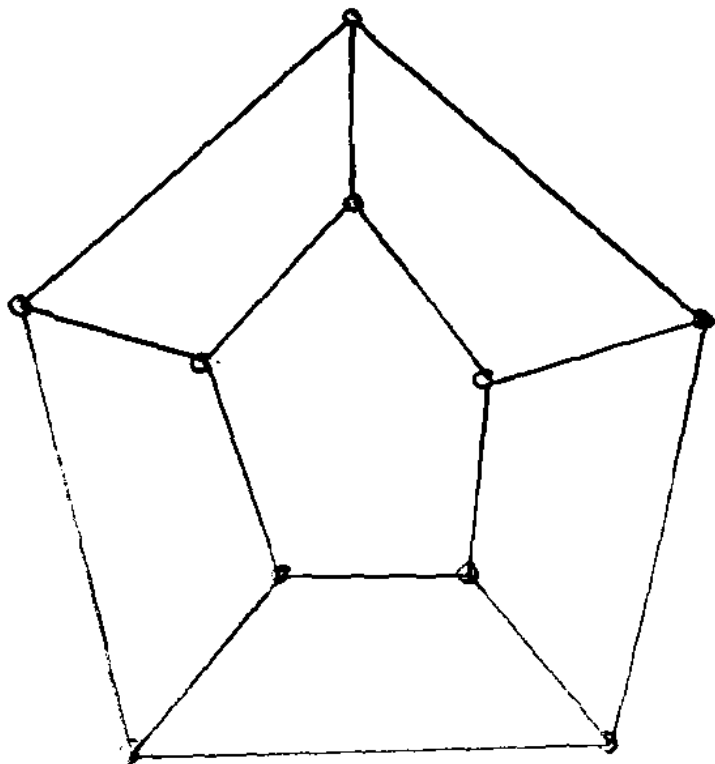
4-edge-conn. is not sufficient



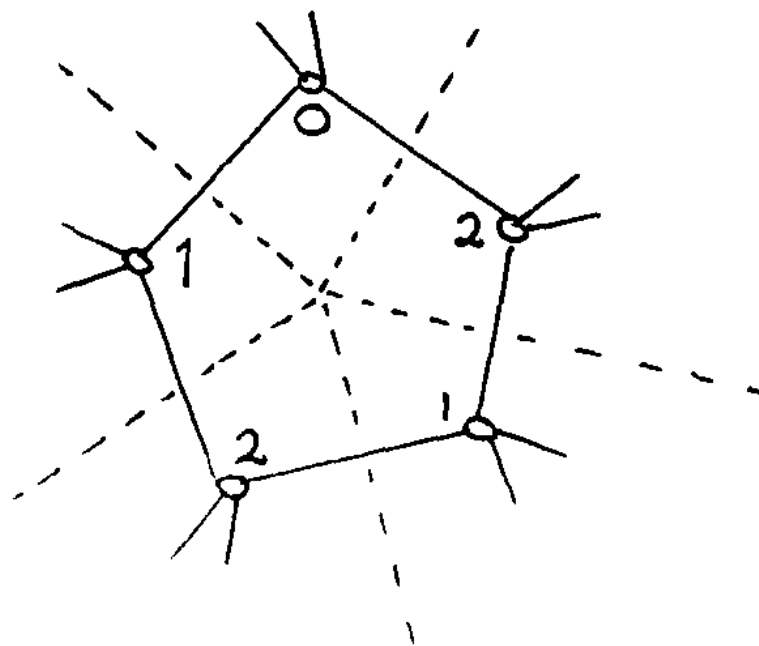
Grötzsch's theorem (1959)

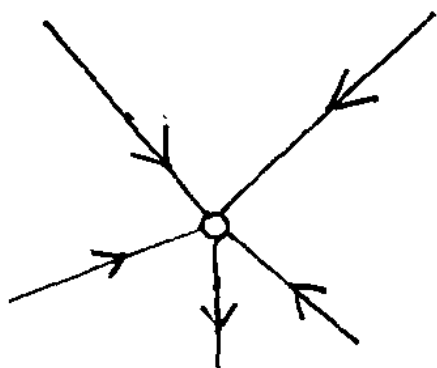
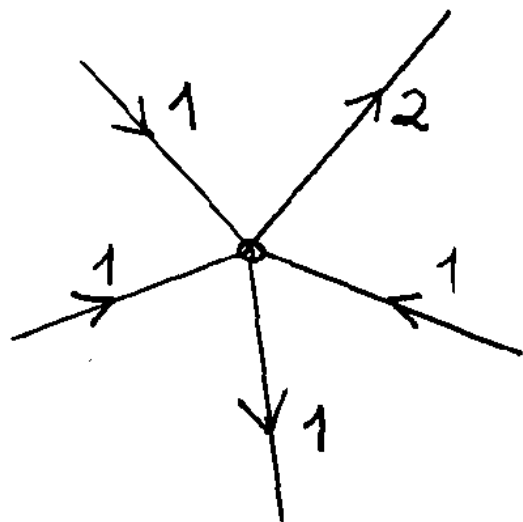
G planar, triangle-free graph \implies

$$\chi(G) \leq 3$$



3-coloring a planar graph $G \implies$
 nowhere ^{zero} 3-flow in G^* \implies
 balanced mod 3 orientation in G^*





Every vertex is balanced
mod 3

Tutte orientation

Grötzsch's theorem

Every planar multigraph
without 1-edge-cuts
and 3-edge-cuts



has a Tutte orientation.

Conjecture of Tutte (1970?)

Every 4-edge-connected
multigraph has a
Tutte orientation.

Weaker conjecture of
Jaeger (1988) :

4 \rightarrow larger const.
 k_F

Conjecture (JB+CT)

There exist constants

$$k_c, k_g$$

such that

(1) every k_g -edge-conn. multigraph admits all generalized Tutte-orientations.

(2) every k_c -edge-conn. graph has a claw-decomposition (provided its size is $\equiv 0 \pmod{3}$).

26a

Generalized Tutte-orientations of a graph G :

For every vertex v , let $q(v) \in \{0, 1, 2\}$.

Assume also

$$\sum_{v \in V(G)} q(v) \equiv |E(G)| \pmod{3}$$

Find an orientation of G such that

$$d^+(v) \equiv q(v) \pmod{3} \quad \forall v \in V(G).$$

Theorem (JB+CT)

If one of k_t, k_c, k_g exists, then they all exist.

Moreover, $k_c \leq k_g$,
 $k_t \leq k_c + 5$
 $k_g \leq 2k_t + 2$

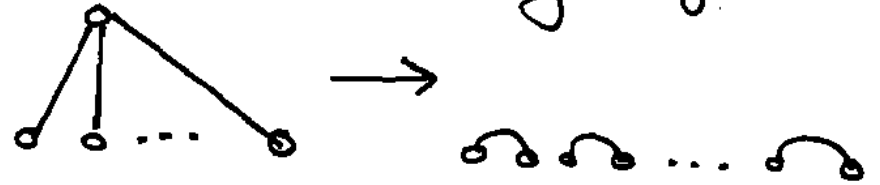
Assume every 10-edge-conn. graph has a claw-decomposition (if its size is $\equiv 0 \pmod{3}$)

Then every 14-edge-conn. multigraph has a Tutte-orientation.

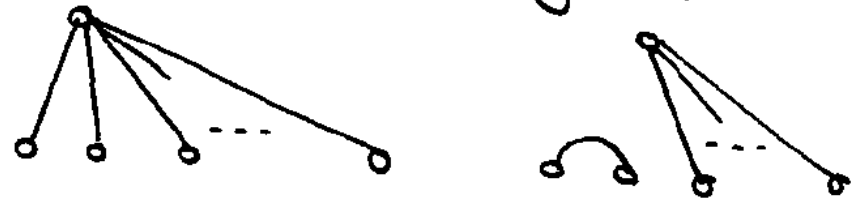
Proof (by induction)

Case 1:  Contract.

Case 2: Vertex of degree 14



Case 3: Vertex of degree ≥ 16



Case 4: 15-regular graph.
Use the claw-decomposition



Conjecture (Jaeger 1988)

For every odd natural number k , there exists a natural number $j(k)$ such that every $j(k)$ -edge-connected multigraph has an orientation which is balanced modulo k .

Equivalent to:

Every $j'(k)$ -edge-conn. graph has a

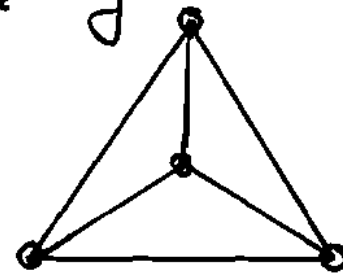
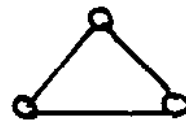
$K_{1,k}$ -decomposition



Thm (Barat + CT)

Let G be a triangulation of a surface.

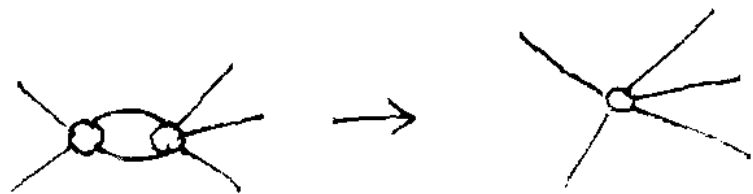
Then G admits all generalized Tutte orientations unless G is one of



Proof in the weakening where the triangulation has at least one double edge.

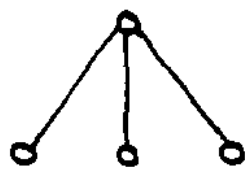
More generally, if a graph G has at least one multiple edge and the contraction of any collection of edges results in multiple edges, then G admits all generalized Tutte-orientations

Proof: Contract a multiple edge and use induction



Stronger (and dangerous)
Conjecture (JB + CT)

For every tree T , there exists a natural number C_T such that every C_T -edge-connected graph G has a T -decomposition



Tutte's conjecture

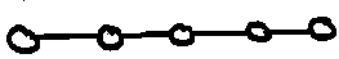


Jaeger's —

What about



?

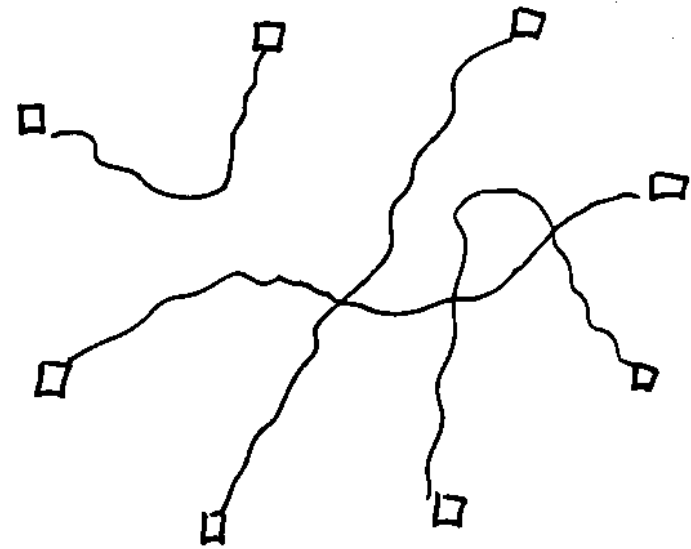
Theorem G has a
 -decomposition
 if $|E(G)| \equiv 0 \pmod 4$
 and G is
 $10^{10} 10^{14}$ - edge-conn.

Trivial case: All vertices
 even degree, $girth \geq 5$.

Lemma G connected
 $S \subseteq V(G)$
 $|S|$ even

Then G has a collection
 of edge-disjoint paths
 from S to S such
 that each vertex of S
 is the end of precisely
 one path.

S : 



A complete collection
of S-paths

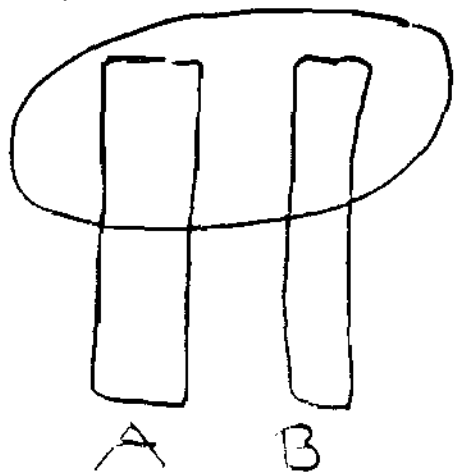
Nearly all can be
 chosen to be even

Lemma

G $(2k-1)$ -edge-connected.

Then $V(G)$ has a bipartition $V(G) = A \cup B$ such that $G[A, B]$ is k -edge-connected.

Proof Select A, B such that $G[A, B]$ has as many edges as possible.

Lemma. G 7-edge-conn.

$S \subseteq V(G)$

$|S|$ even.

Then G has a complete collection of S -paths such that all - except possibly one - is of even length.

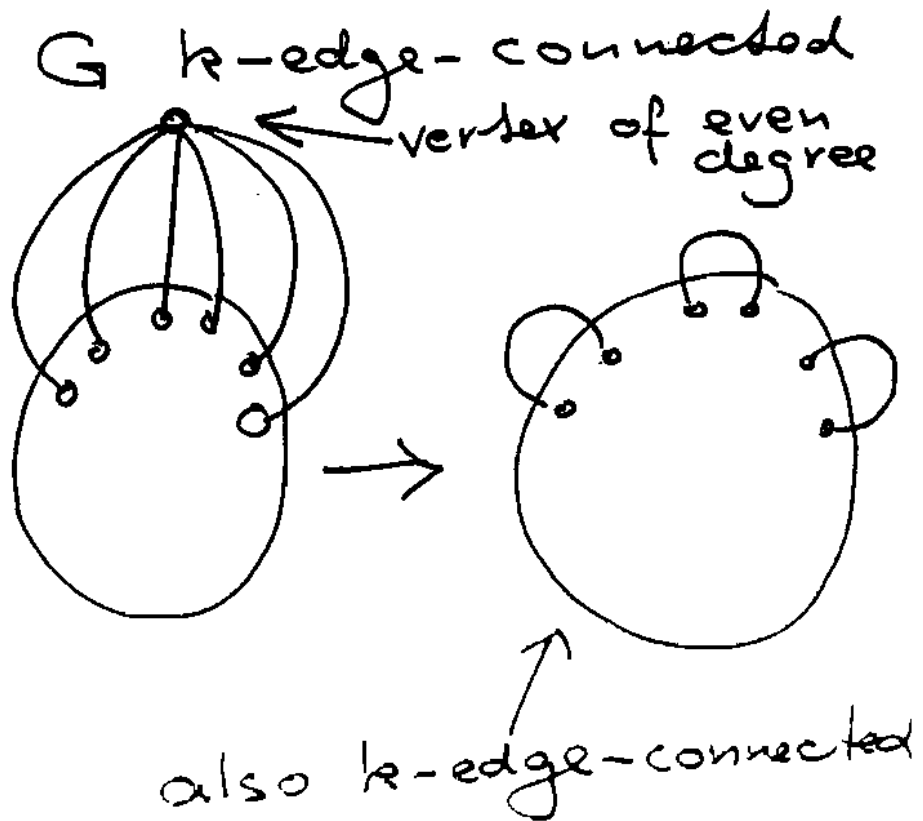
Proof. Choose $V(G) = A \cup B$ such that $G[A, B]$ is 4-edge-conn.

Let T_1, T_2 be edge-disj. spanning trees.

$S_1 = S \cap A, S_2 = S \cap B$.

Take a complete collection of S_i -paths in $T_i, i=1, 2$.

Special case -
 due to Lovász -
 of Mader's lifting
theorem



Thm. $m = 2^r$
 G δ^m -edge-conn.
 girth $> m$.

Let $S \subseteq V(G)$, $|S|$ even

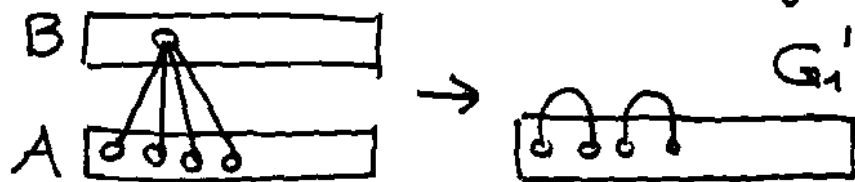
Then G has a
 complete collection of
 S -paths such that
 almost all have
 length $\equiv 0 \pmod{m}$

Proof Let $V(G) = A \cup B$
 such that $G[A, B]$
 has large edge-connecti-
 vity

$$G[A, B] = G_1 \cup G_2$$

G_1, G_2 spanning and of large edge-conn.

Delete edges in G_1 such that all vertices in B have even degree



Mader's lifting

Apply induction of G_1' with $S \cap A$ instead of S .

Do the same for A and G_2

Thm. $m = 2^f$

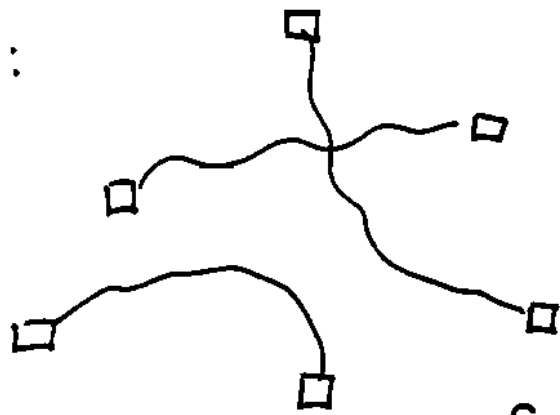
G 8^{m+1} -edge-conn.
girth $> m$
size divisible by m .

Then G has an m -path-decomposition.

Proof Write $G = G_1 \cup G_2$ where G_i is spanning and of large edge-conn.
 $S =$ the odd vertices in G .

Apply previous thm. to G_1, S .

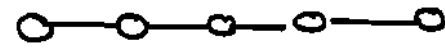
In G_1 :



Paths connecting S
of length divisible by m .
Delete the edges and
decompose into m -paths.

Because of G_2 the
remaining graph
is highly connected
and (almost)
Eulerian

4-path-decomposition



Thm If G is
207-edge-connected
then G has a set
 E of ≤ 6 edges s.t.
 $G-E$ has a 4-path-
decomp.

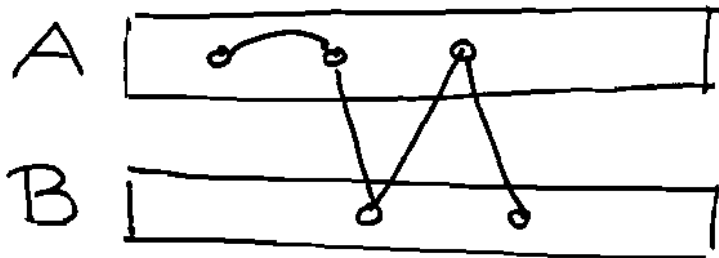
Idea of proof:

$$V(G) = A \cup B$$

$G[A, B]$ is 104-edge-
conn.

and has 52 edge-disj.
spanning trees.

Use a few spanning trees to take care of the edges in $G(A), G(B)$



So we have a bipartite graph.

Then repeat the proof in the large girth case.

The main problem is that Mader's lifting theorem creates multiple edges.

We have now decomposed a 207 -edge-conn. graph into 4 -paths except ≤ 6 edges.

In order to include these 6 edges we raise 207 to $10^{10^{10^{14}}}$.

Conjecture (CT, 1983)

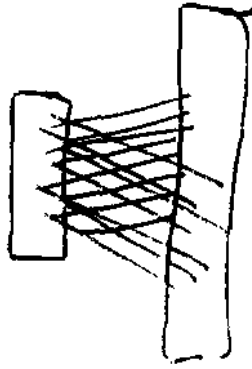
$\forall m, g \exists f(m, g) :$
 $G \text{ } f(m, g)\text{-edge-conn.}$

\Downarrow
 $G \supseteq H \quad m\text{-edge-conn.}$
 $\text{girth} \geq g$

$$f(m, 4) \leq 2m-1$$

because every $(2m-1)$ -edge-conn. graph contains an m -edge-connected bipartite spanning subgraph.

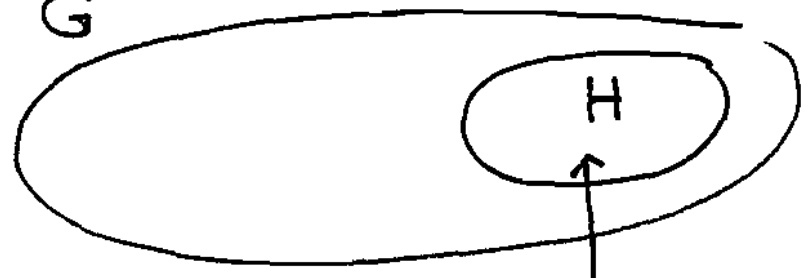
For $g \geq 5$, the subgraph H cannot be spanning.



Kühn, Osthus, 2004

$$f(k, 6) < 64 \cdot (8k)^{3+2 \cdot 11^{64}} 8k^3$$

Assume now G is $10^{10^{14}}$ -edge-conn.



highly conn.
girth > 4

$$H = H_1 \cup H_2 \text{ each spanning highly-conn.}$$

Decompose $G - E(H_1)$ into 4-paths except for 6 edges in H . Use H_2

Conjecture

$$\forall k, g \exists h(k, g):$$

$$G \text{ } h(k, g)\text{-edge-conn.}$$

$$\Downarrow$$

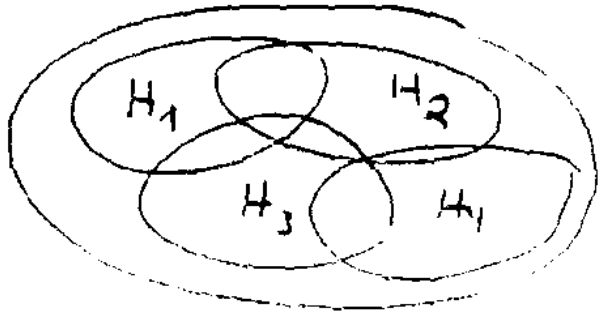
$$G \cong H_1, H_2, \dots, H_m$$

pairwise edge-disjoint
subgraphs of $\text{girth} \geq g$,

edge-conn. $\geq k$

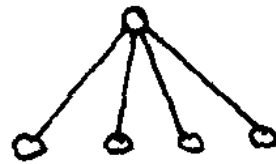
such that

$$V(G) = V(H_1) \cup V(H_2) \cup \dots$$

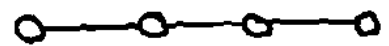


If true, then the tree-decomposition conjecture holds for every path of length a power of 2.

The next natural cases:



4-star



3-path